

Impossibility

DDA Course
week 3

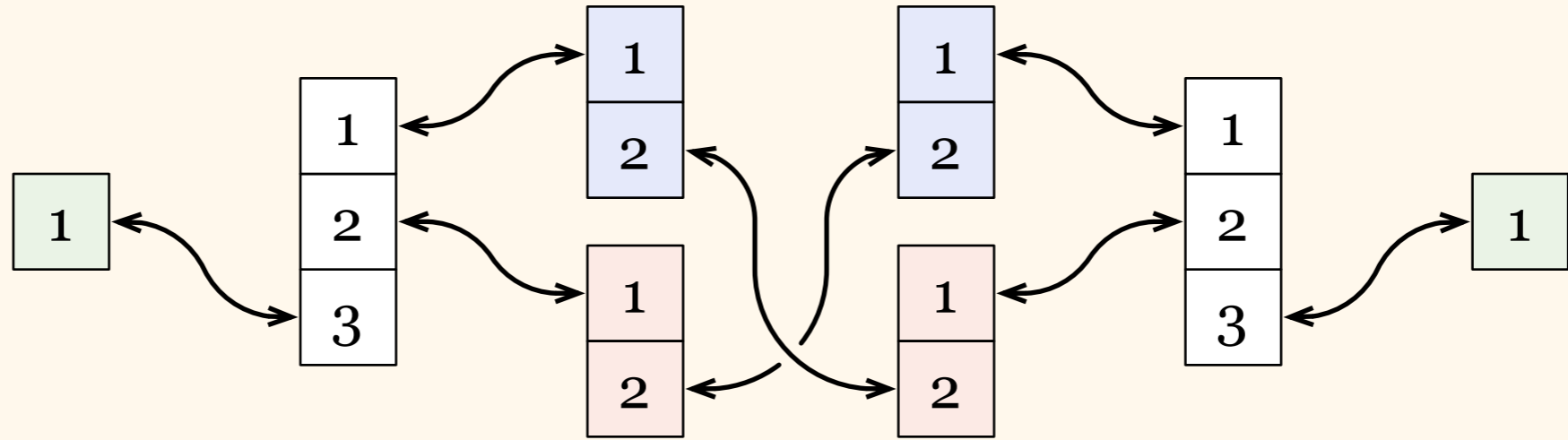
Proof Techniques

- Covering maps
 - problems that cannot be solved at all
- Isomorphic local neighbourhoods
 - problems that cannot be solved quickly

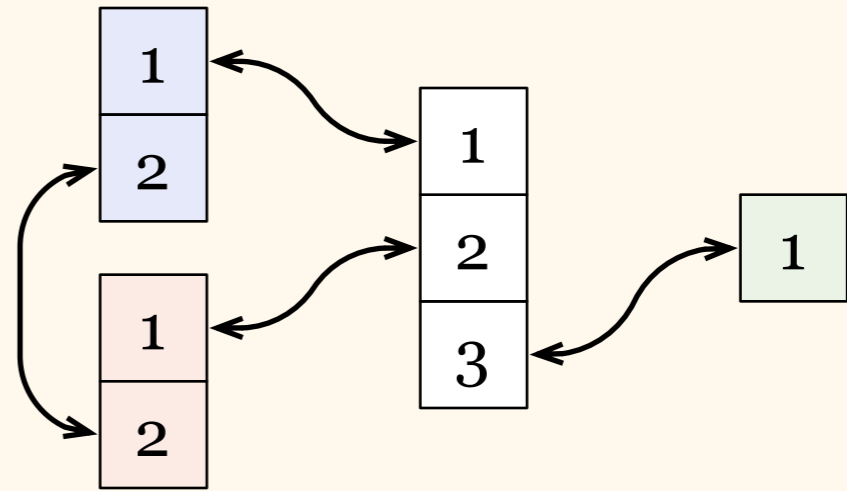
Covering Map

- Networks $N = (V, P, p)$ and $N' = (V', P', p')$
- Surjection $\varphi: V \rightarrow V'$ that *preserves inputs, degrees, connections, and port numbers*

N :



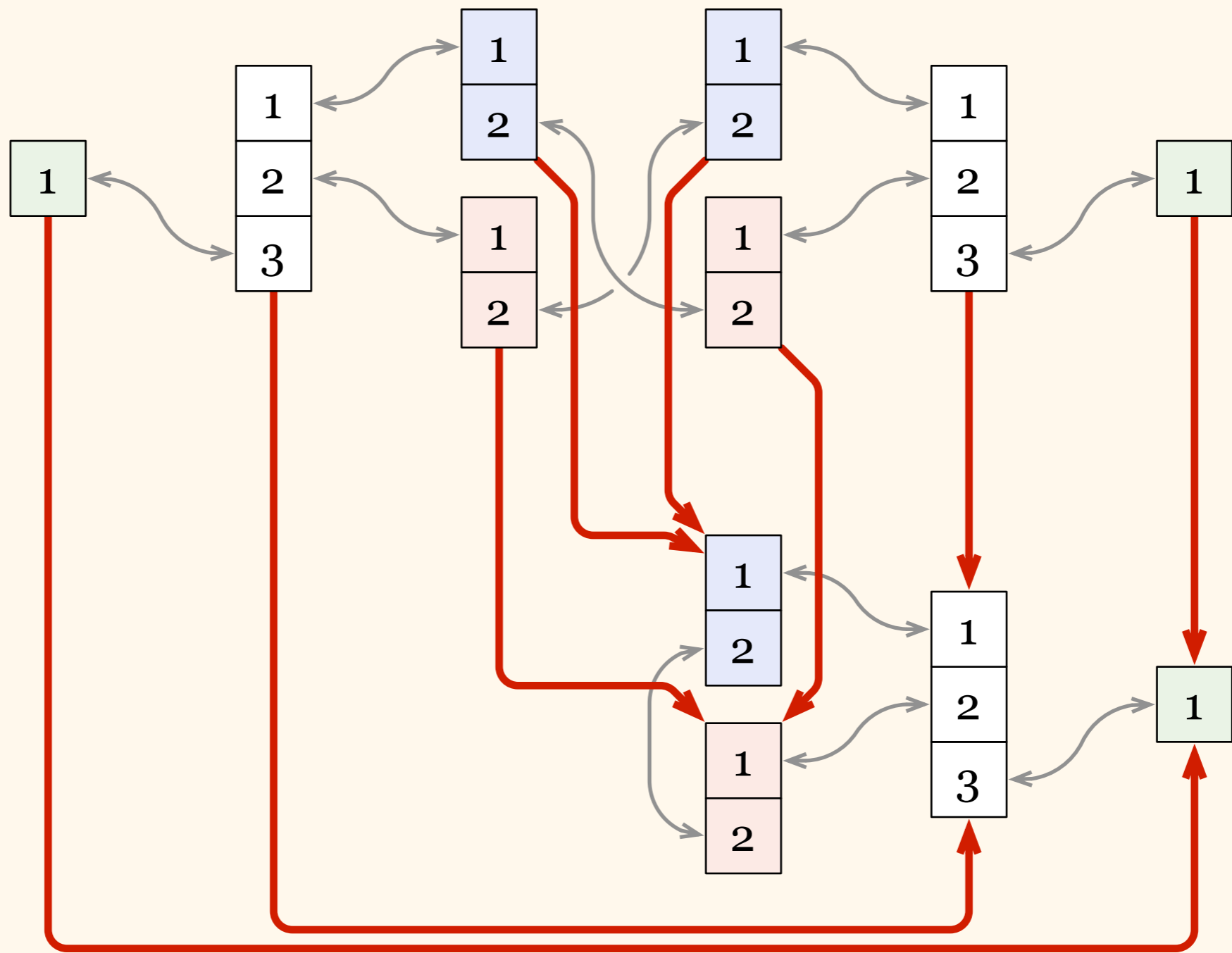
N' :



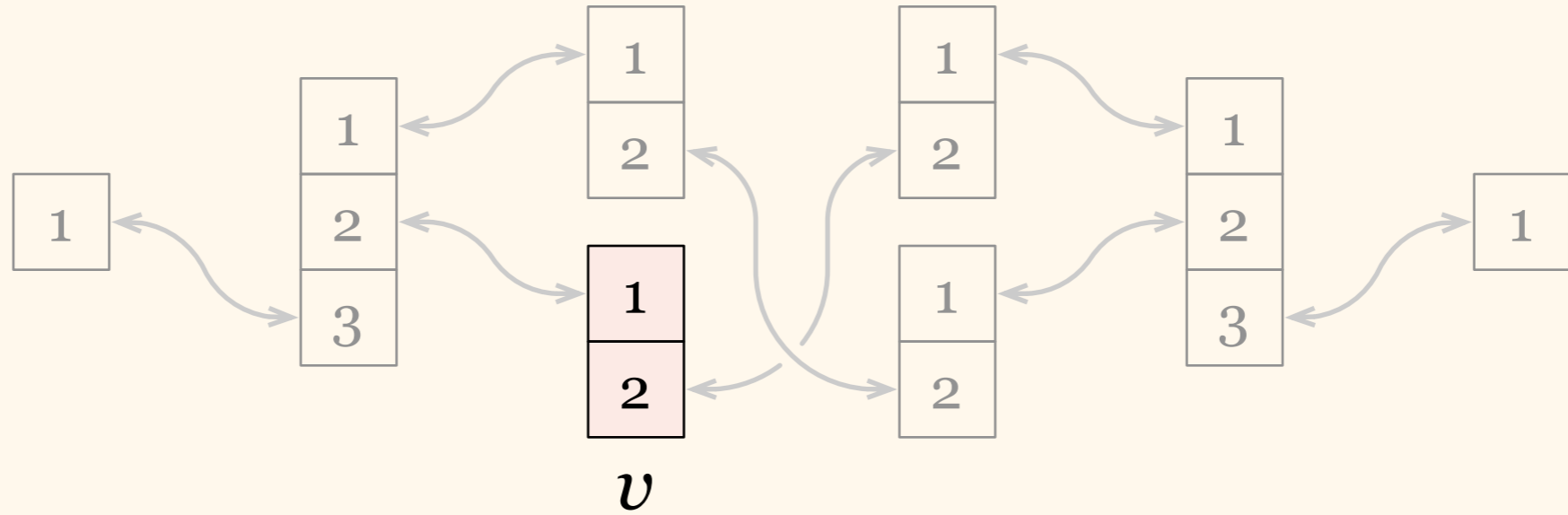
$N:$

φ

$N':$

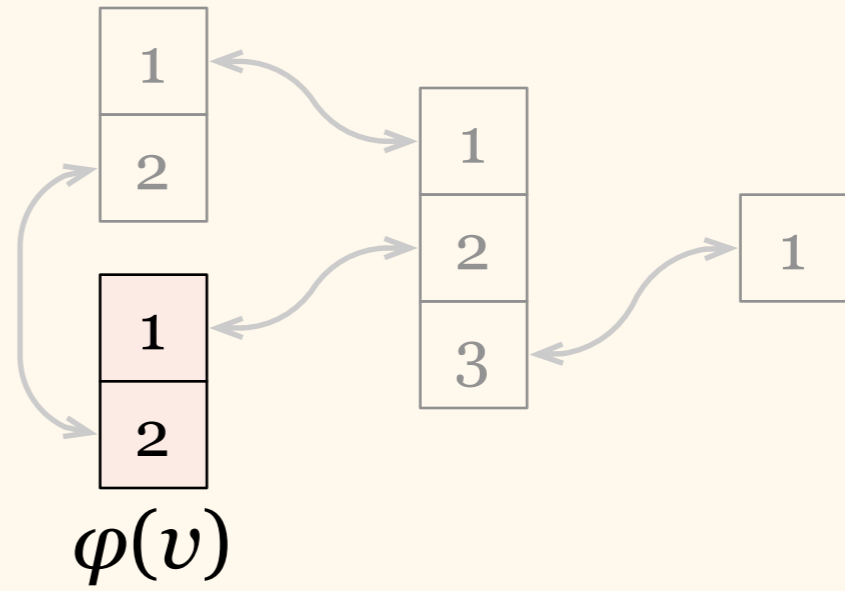


$N:$



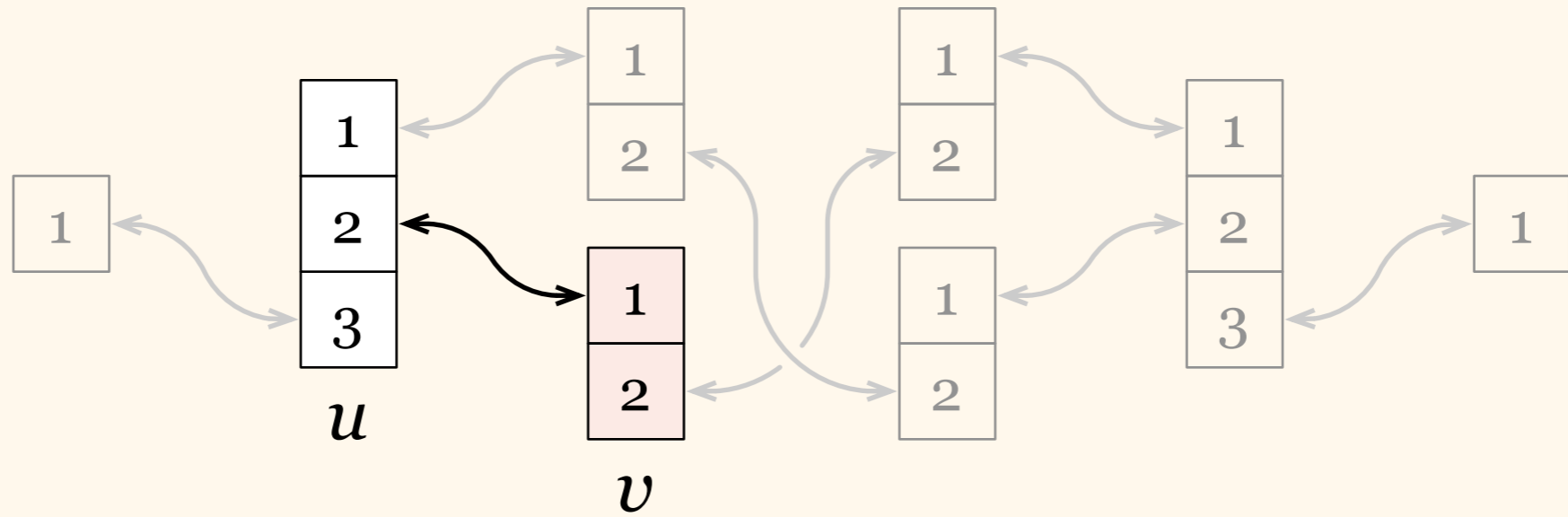
φ

$N':$



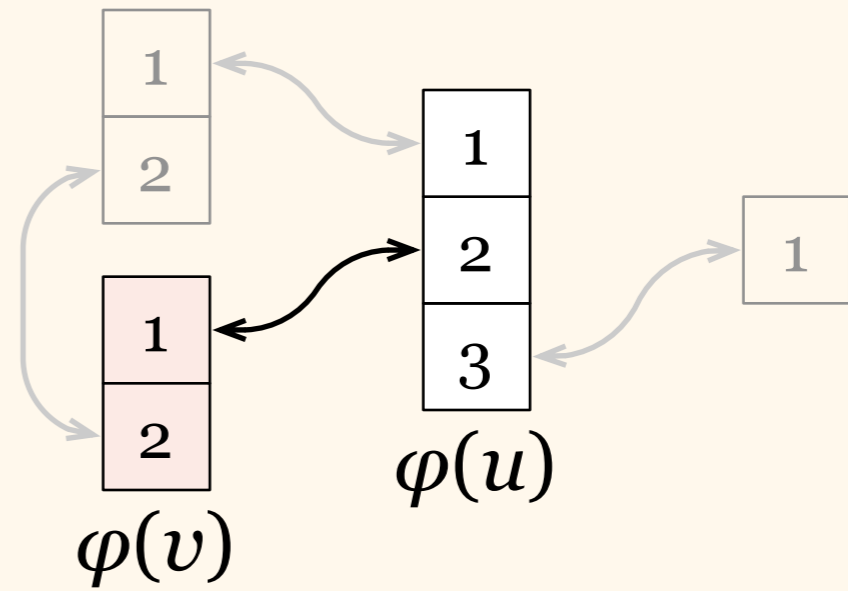
Degrees agree

$N:$

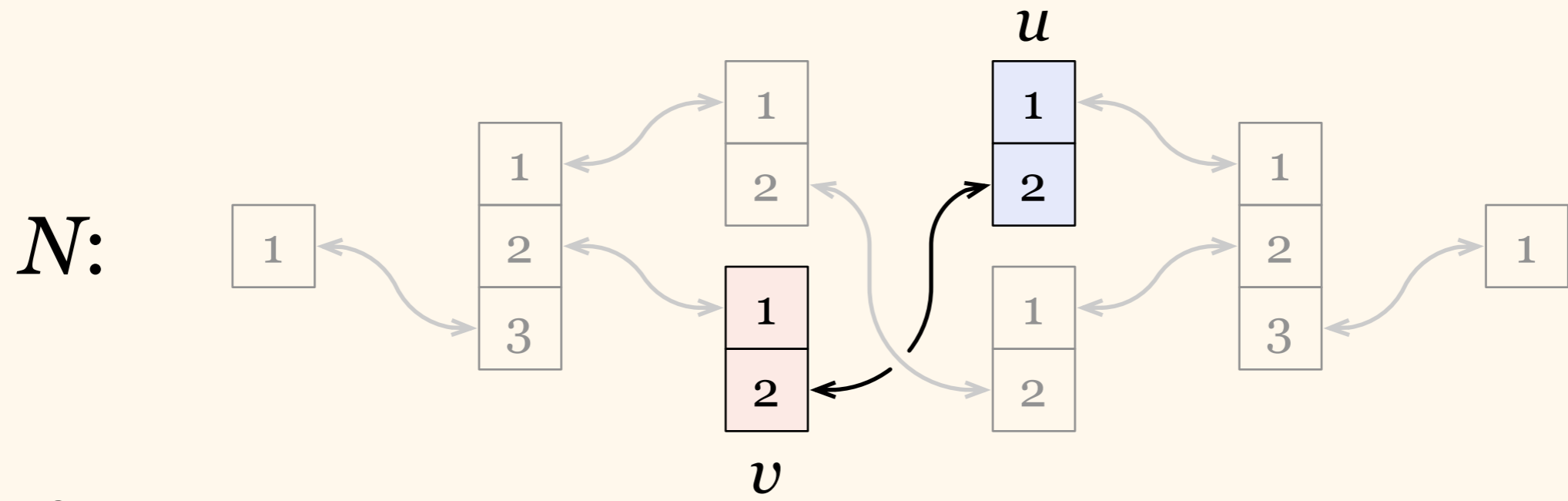


φ

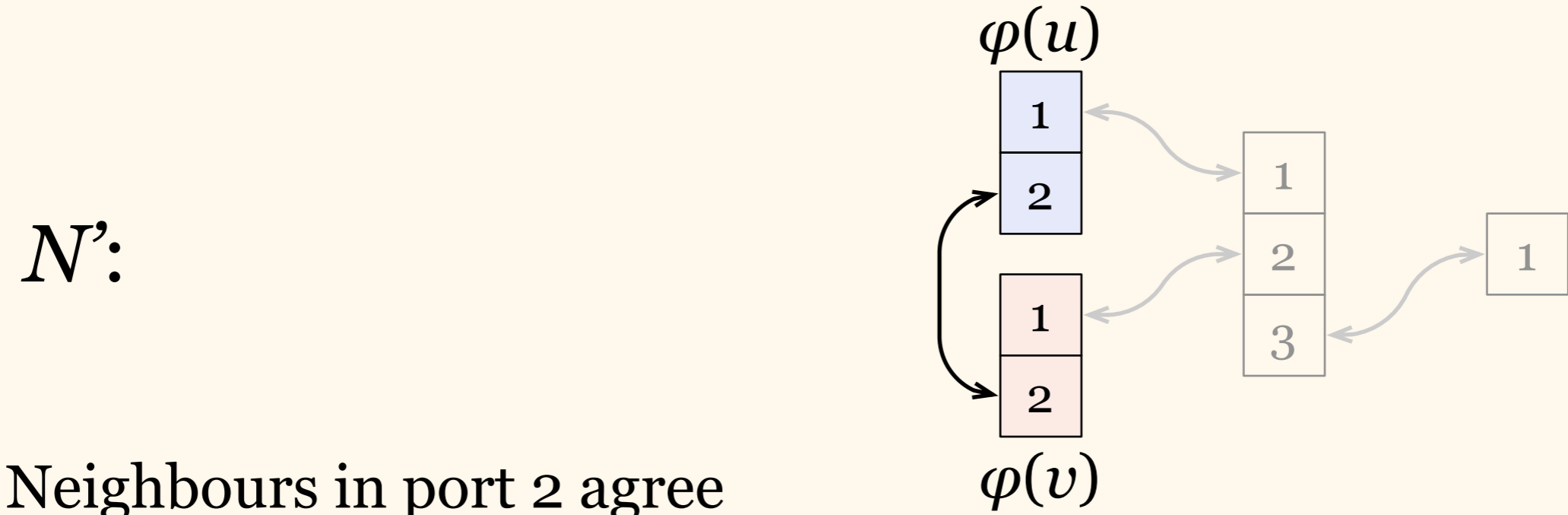
N' :



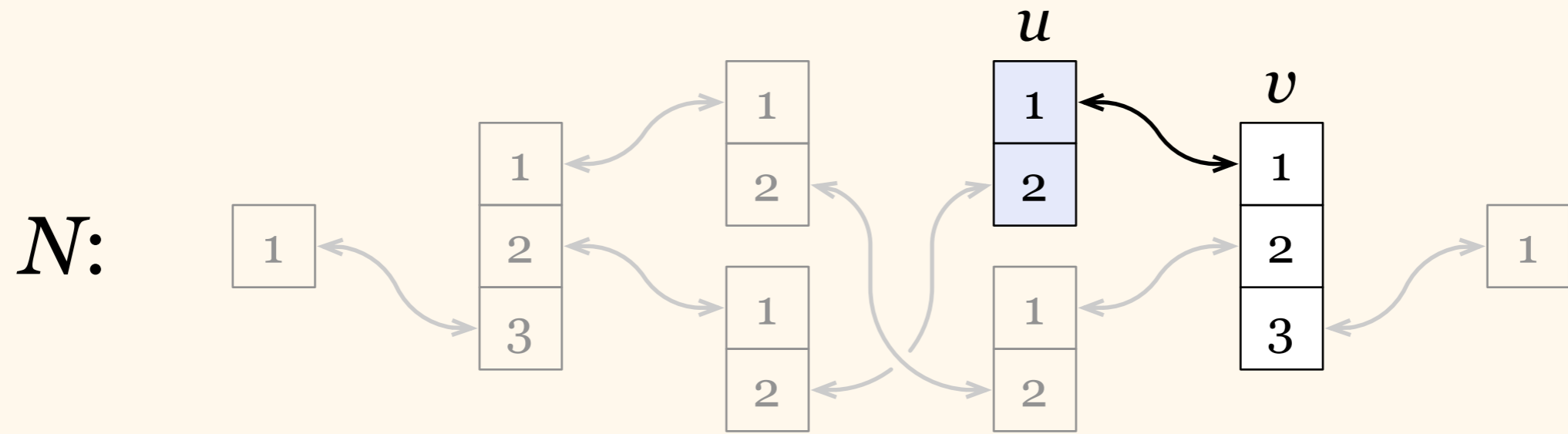
Neighbours in port 1 agree



φ

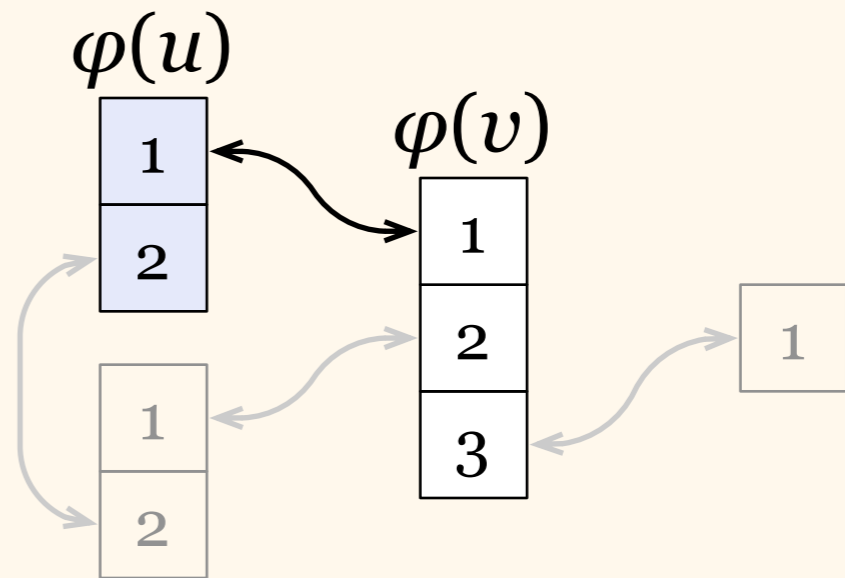


Neighbours in port 2 agree



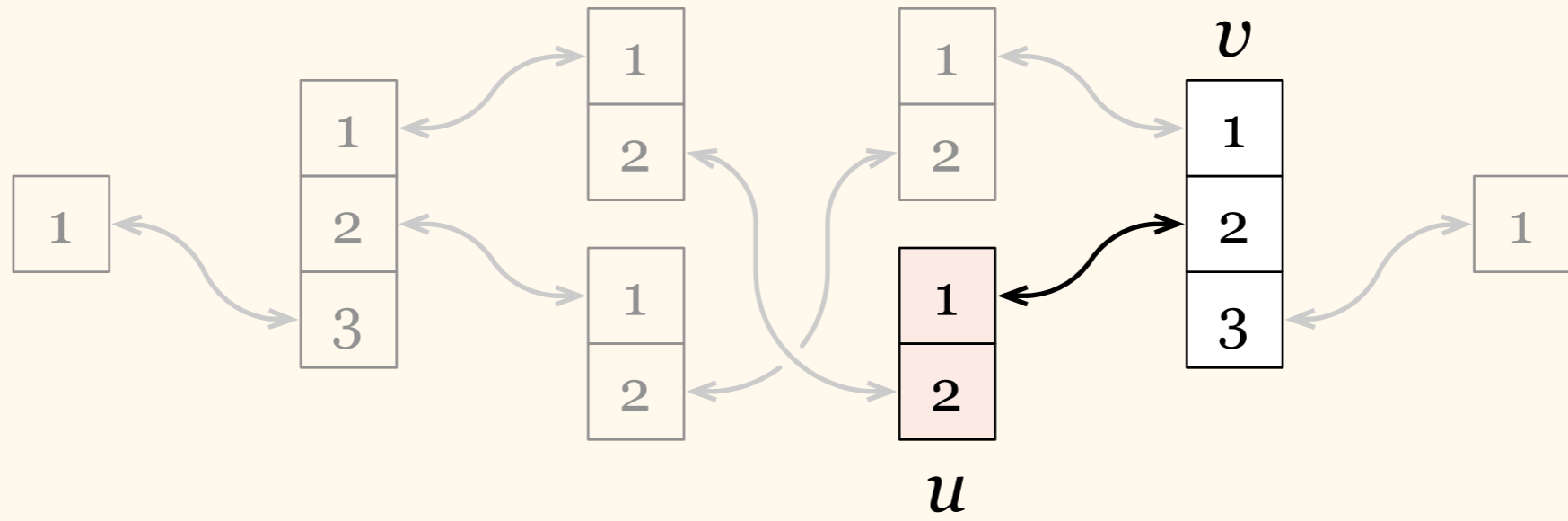
φ

$N':$



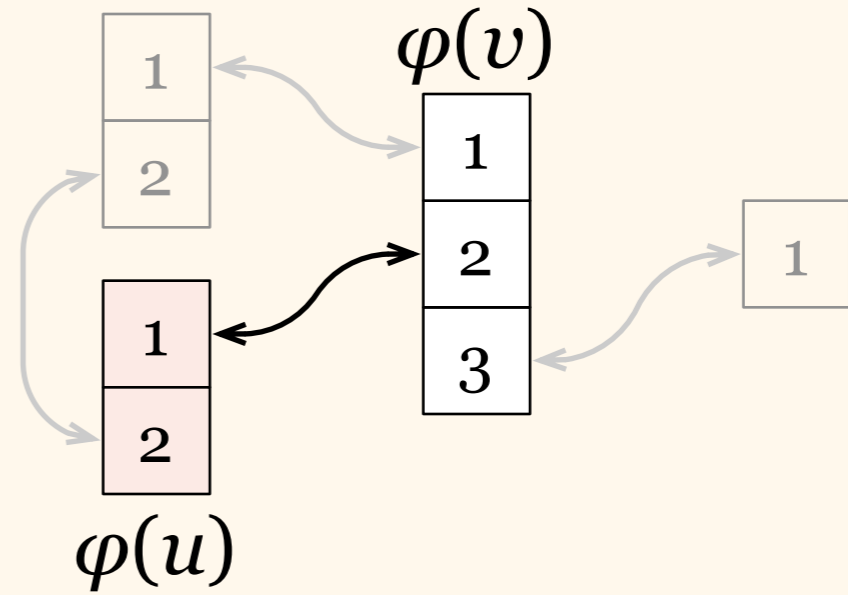
Holds for any pair of nodes

$N:$



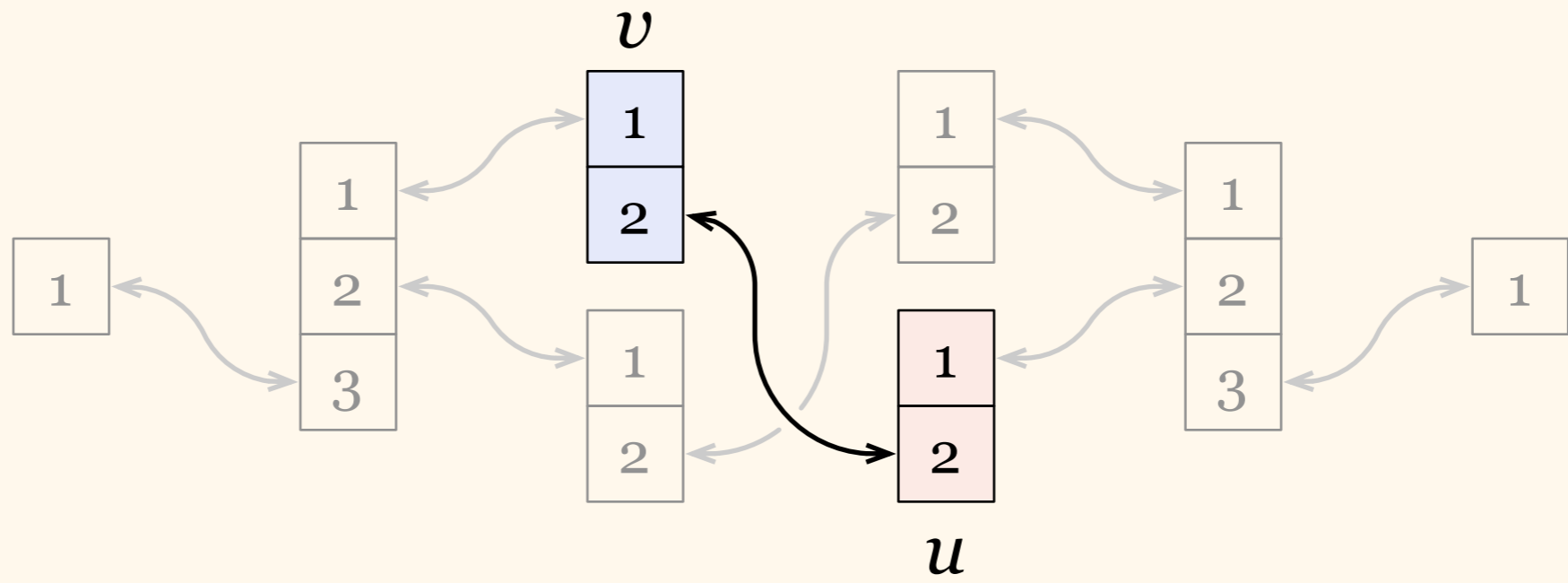
φ

$N':$



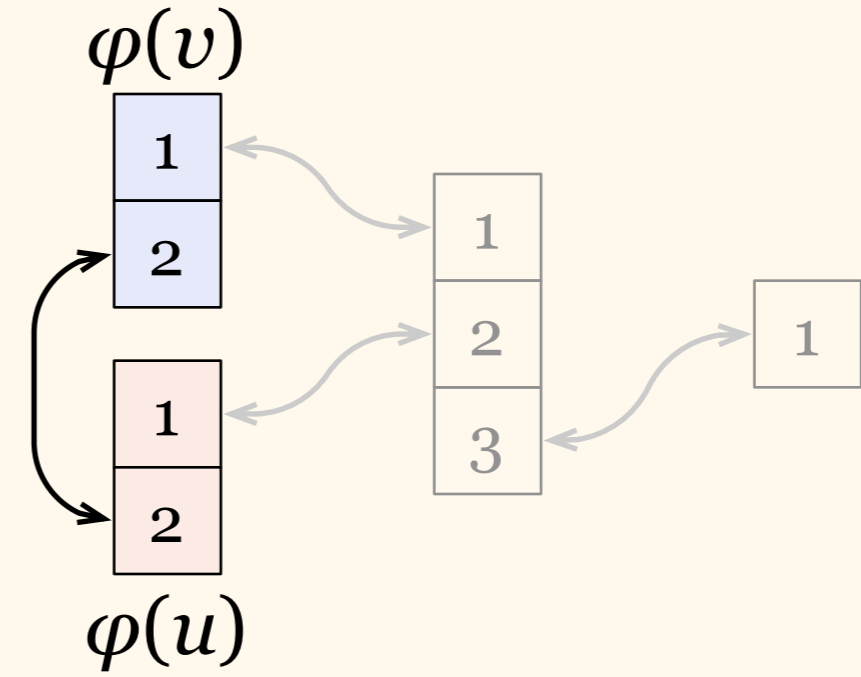
Holds for any pair of nodes

$N:$

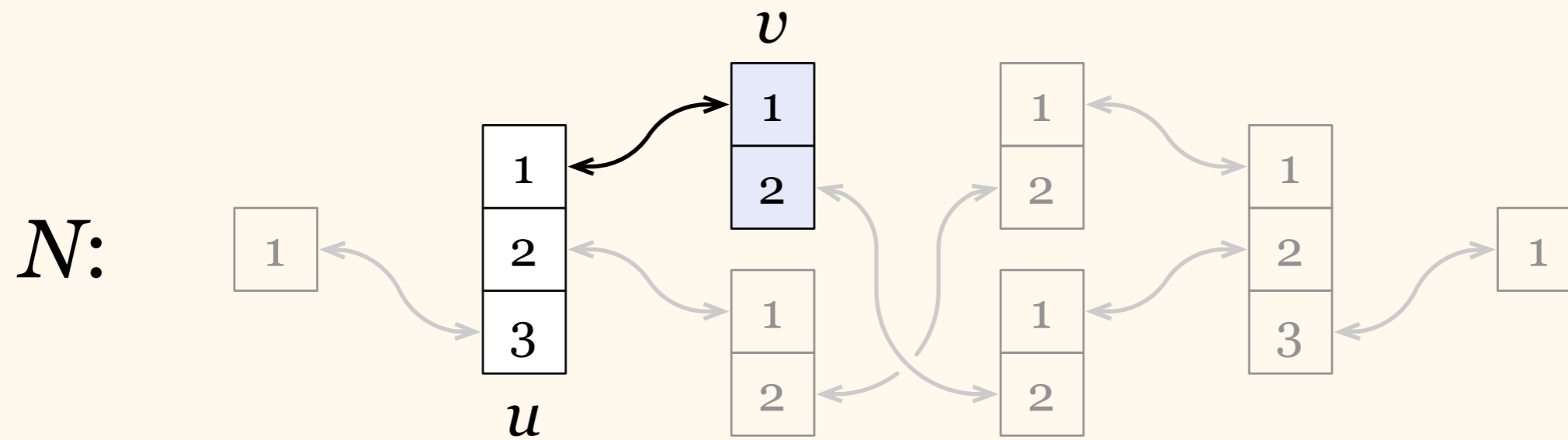


φ

$N':$

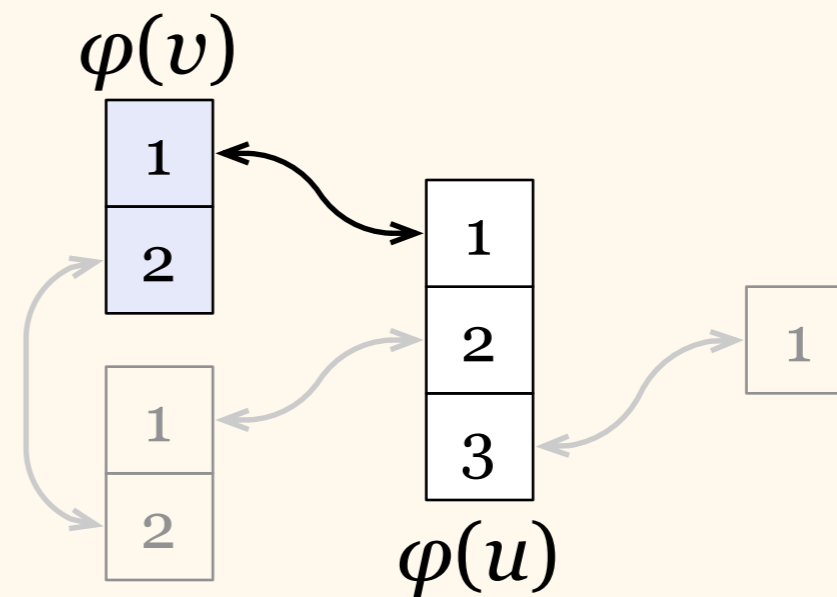


Holds for any pair of nodes



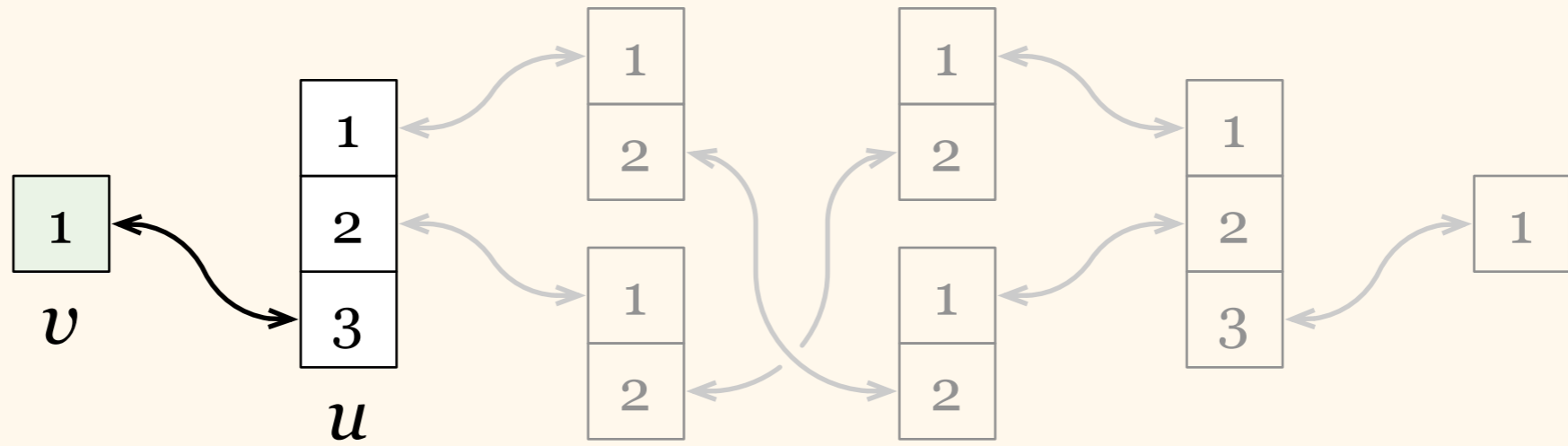
φ

N' :



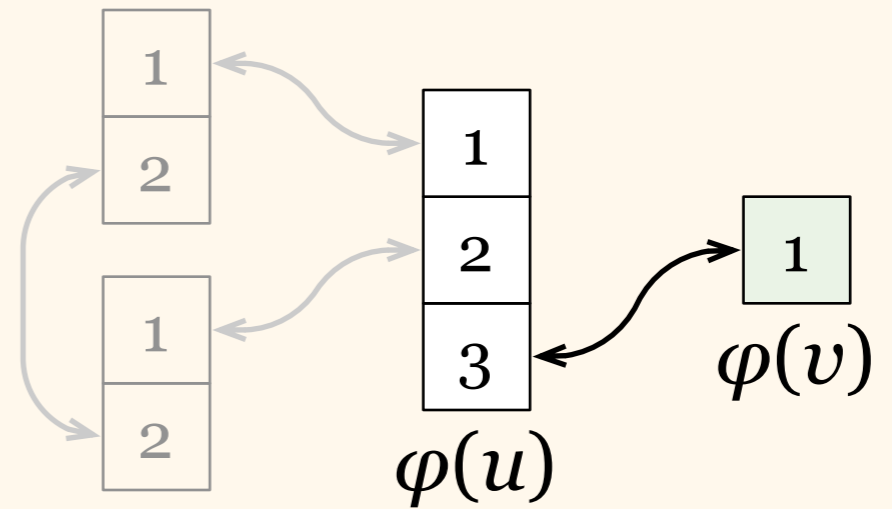
Holds for any pair of nodes

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φ

$N':$



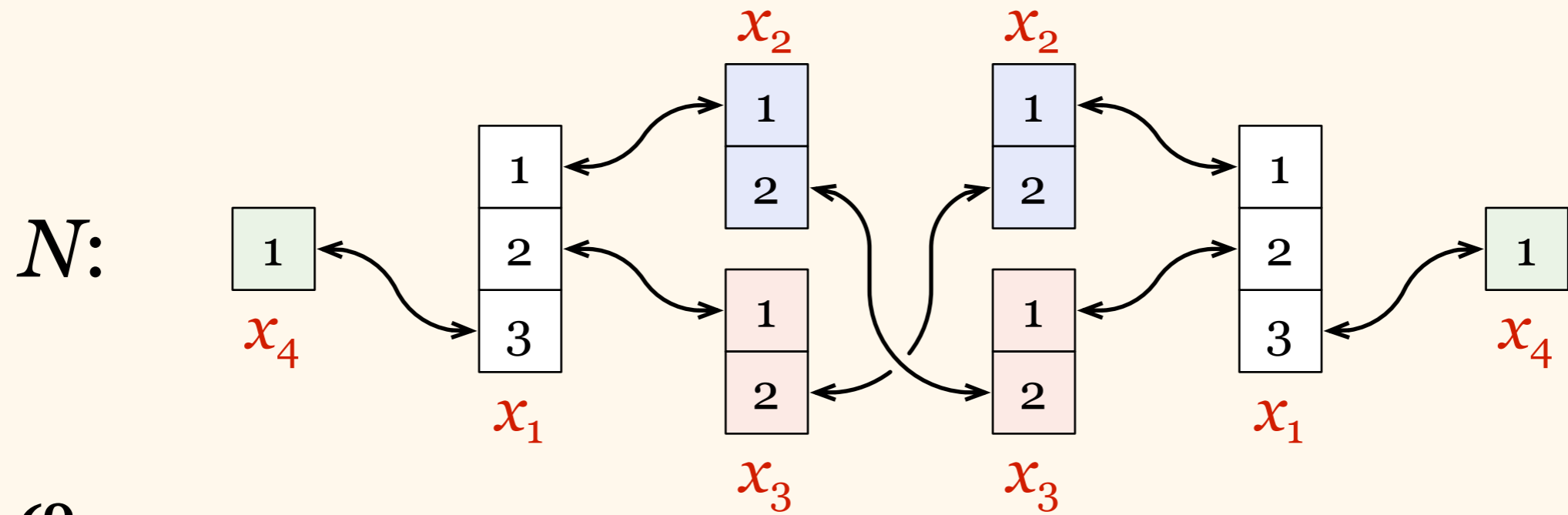
Holds for any pair of nodes

Covering Map

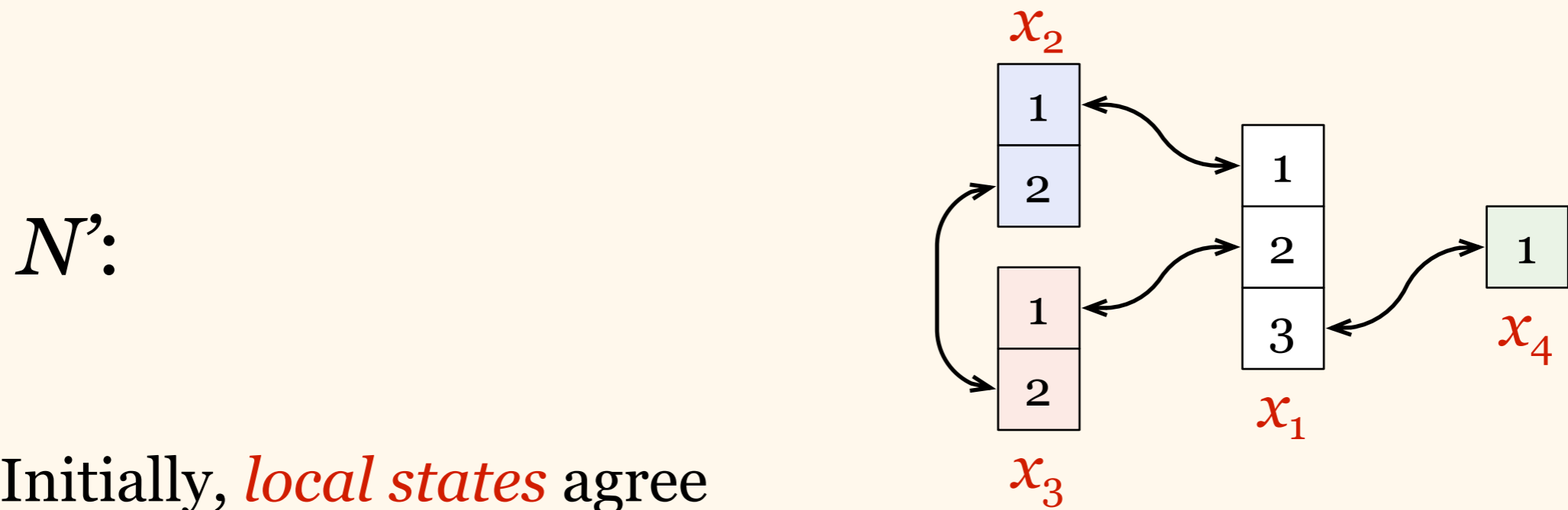
- Networks $N = (V, P, p)$ and $N' = (V', P', p')$
- Surjection $\varphi: V \rightarrow V'$ that preserves inputs, degrees, connections, and port numbers
- **Theorem:** If we run an algorithm A in N and N' , then nodes v and $\varphi(v)$ are *always in the same state*

Covering Map

- **Theorem:** If we run an algorithm A in N and N' , then nodes v and $\varphi(v)$ are *always in the same state*
- **Proof:** By induction
 - before round i : map φ preserves local states
 - during round i : map φ preserves messages
 - after round i : map φ preserves local states

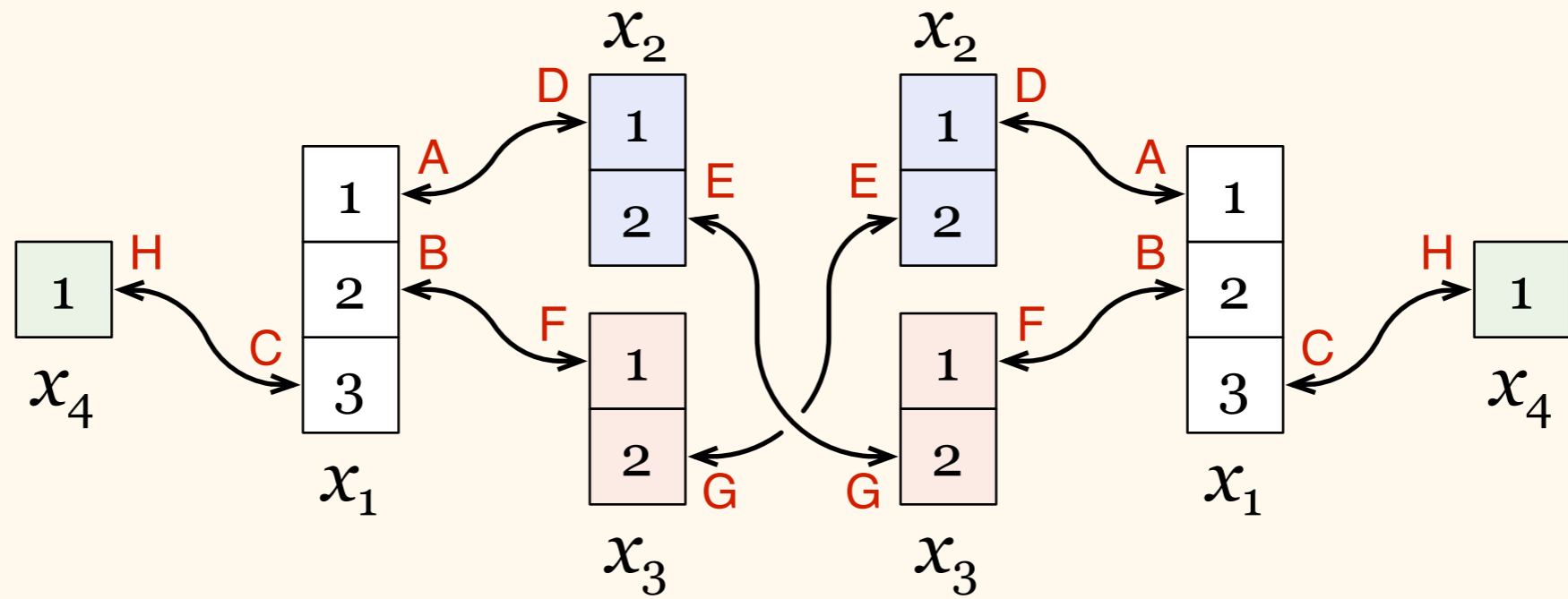


φ



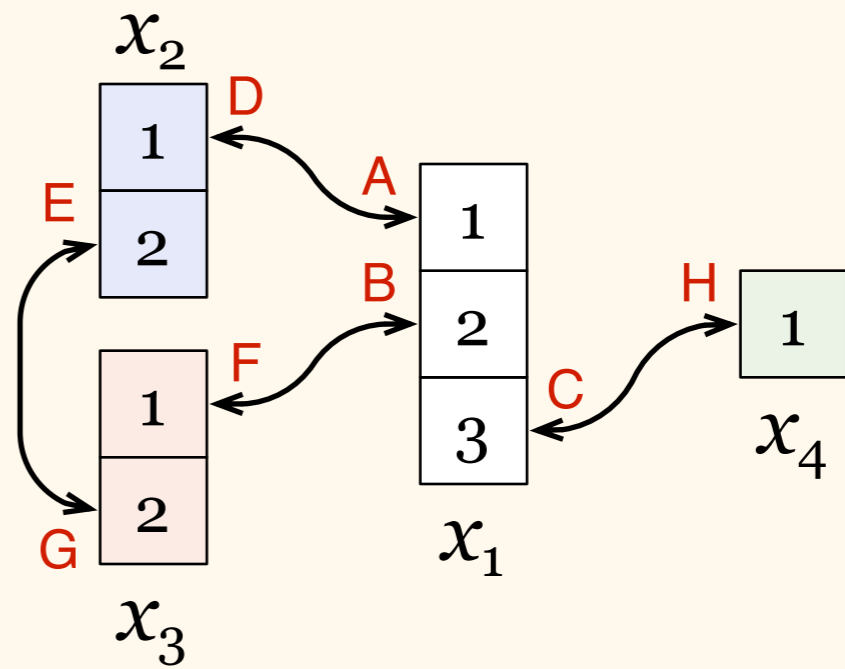
Initially, *local states* agree

$N:$



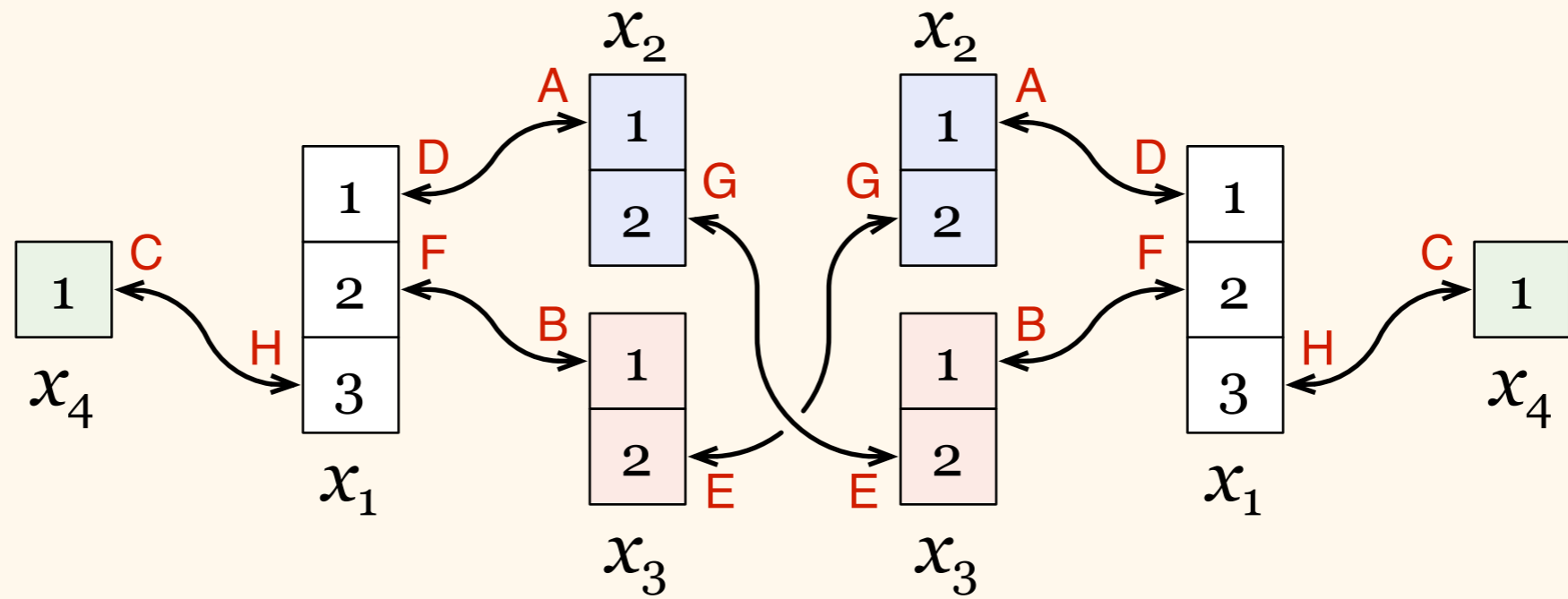
φ

$N':$



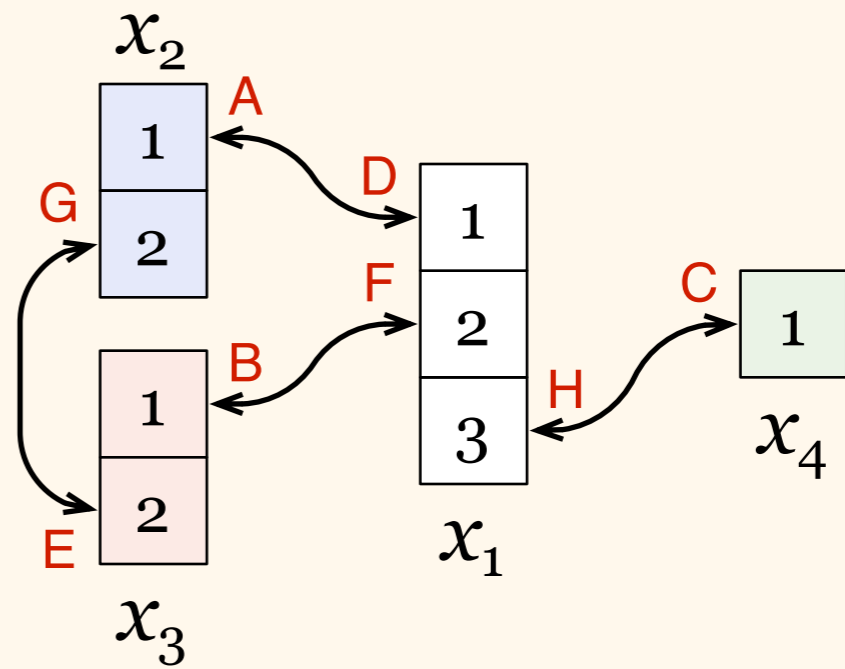
Thus *outgoing messages* agree

$N:$



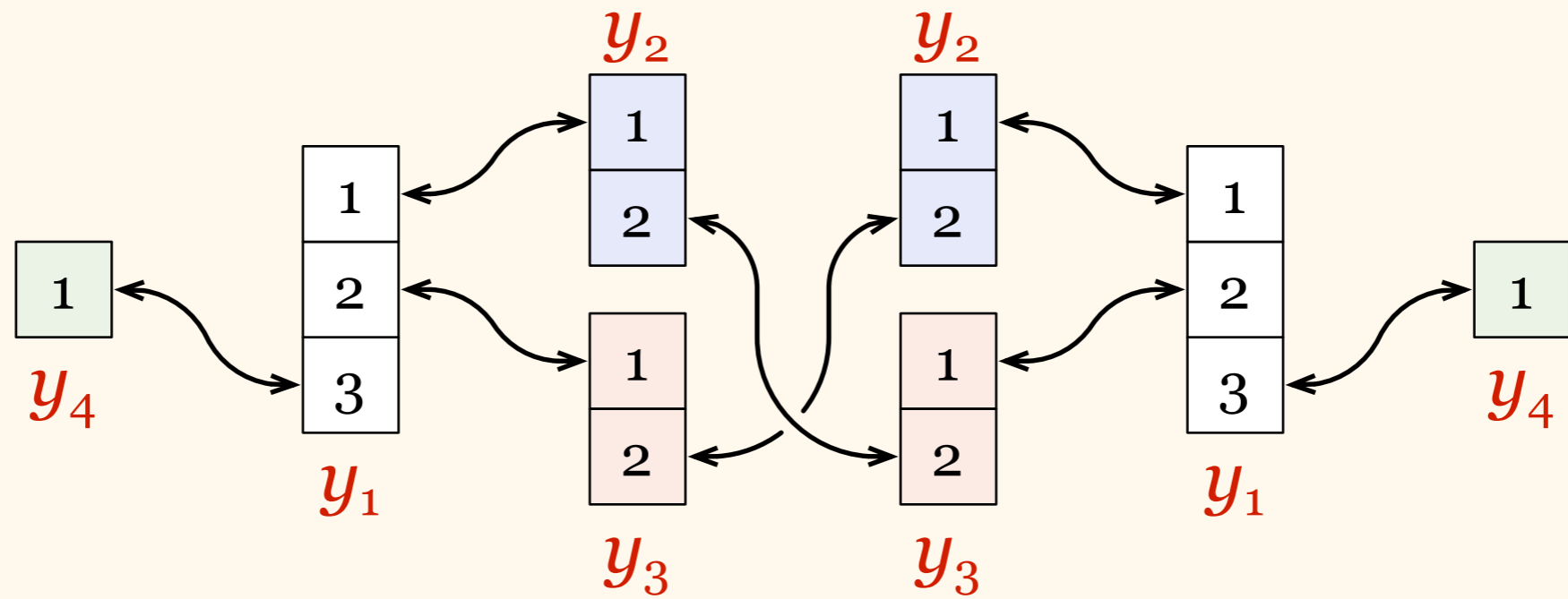
φ

$N':$



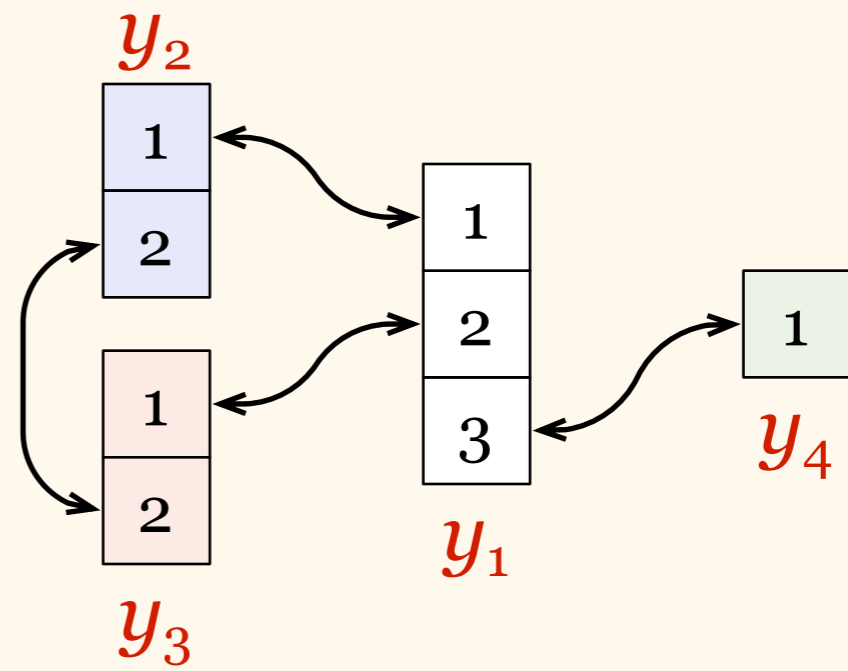
Thus *incoming messages* agree

$N:$



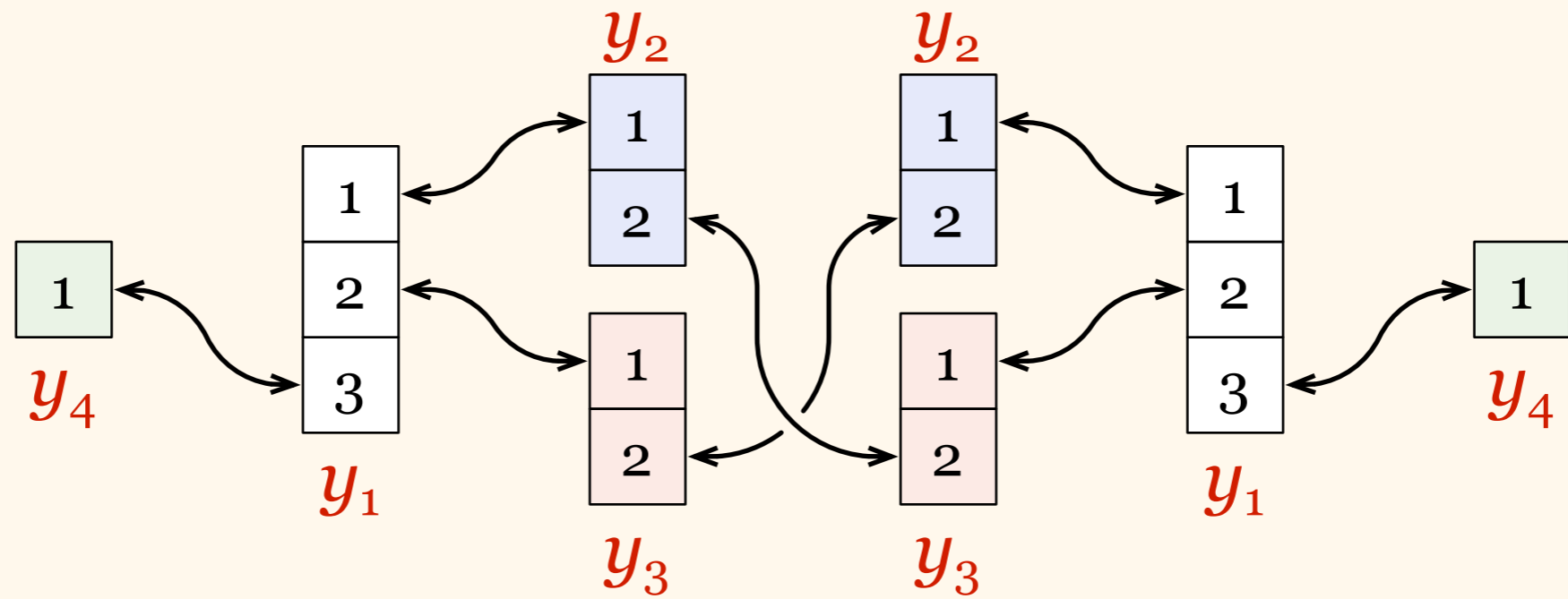
φ

$N':$



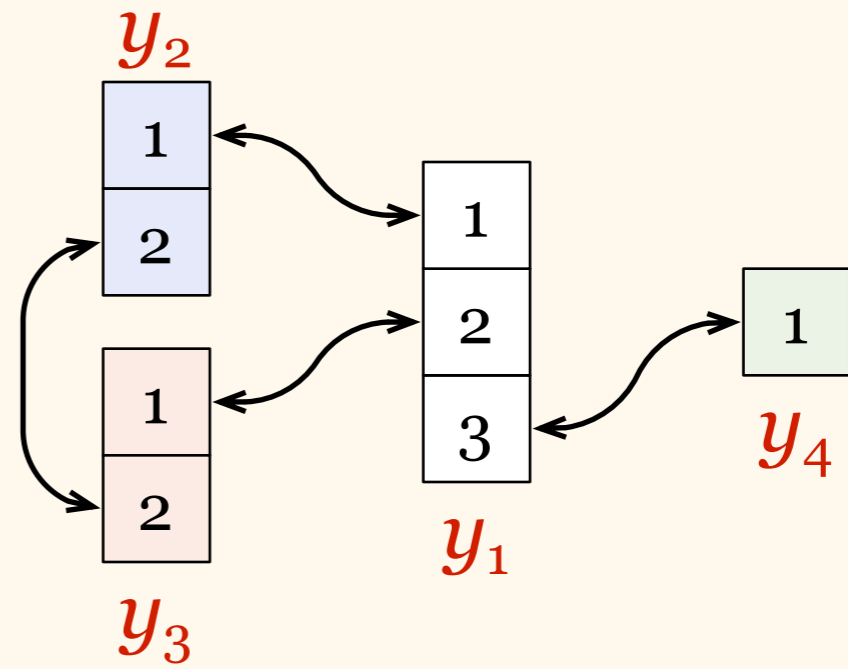
Thus *new local states* agree

$N:$



φ

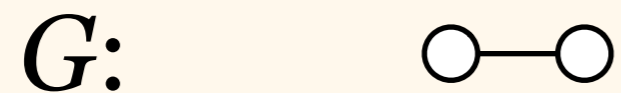
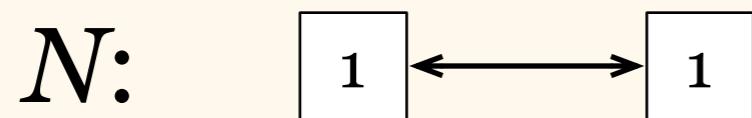
$N':$



By induction, *local outputs* agree

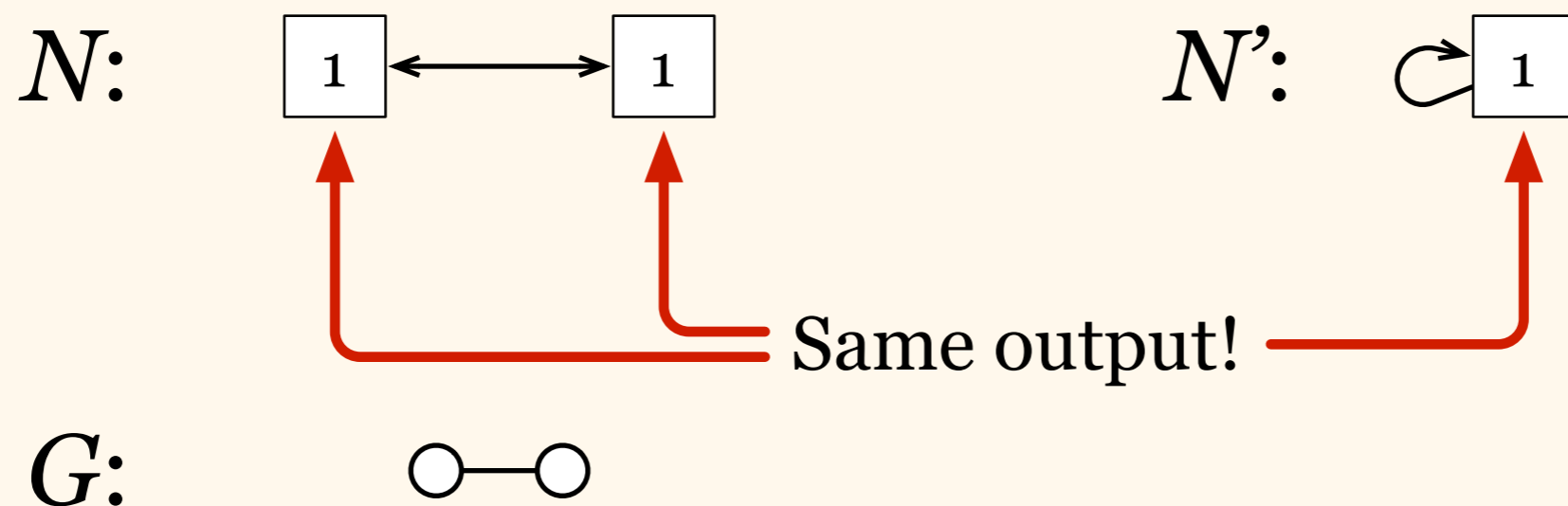
Covering Map

- Application: symmetry breaking in a path graph



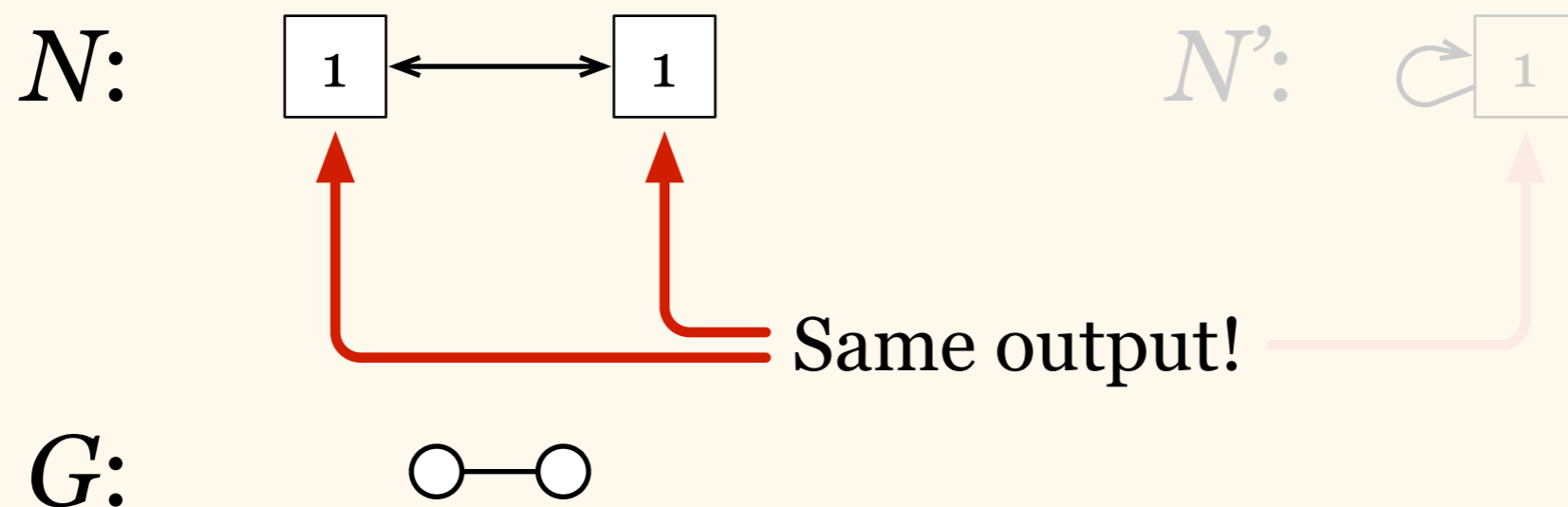
Covering Map

- Application: symmetry breaking in a path graph



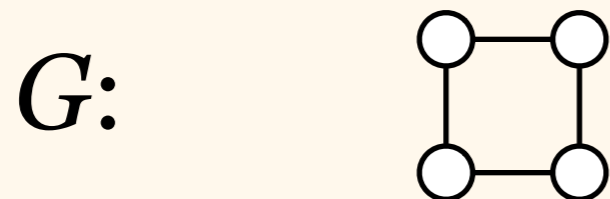
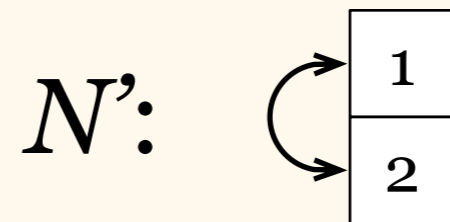
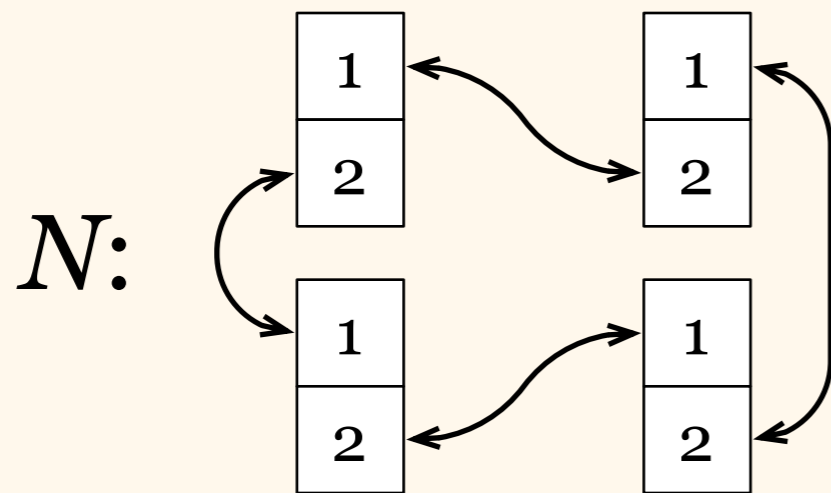
Covering Map

- Application: symmetry breaking in a path graph



Covering Map

- Application: symmetry breaking in a cycle

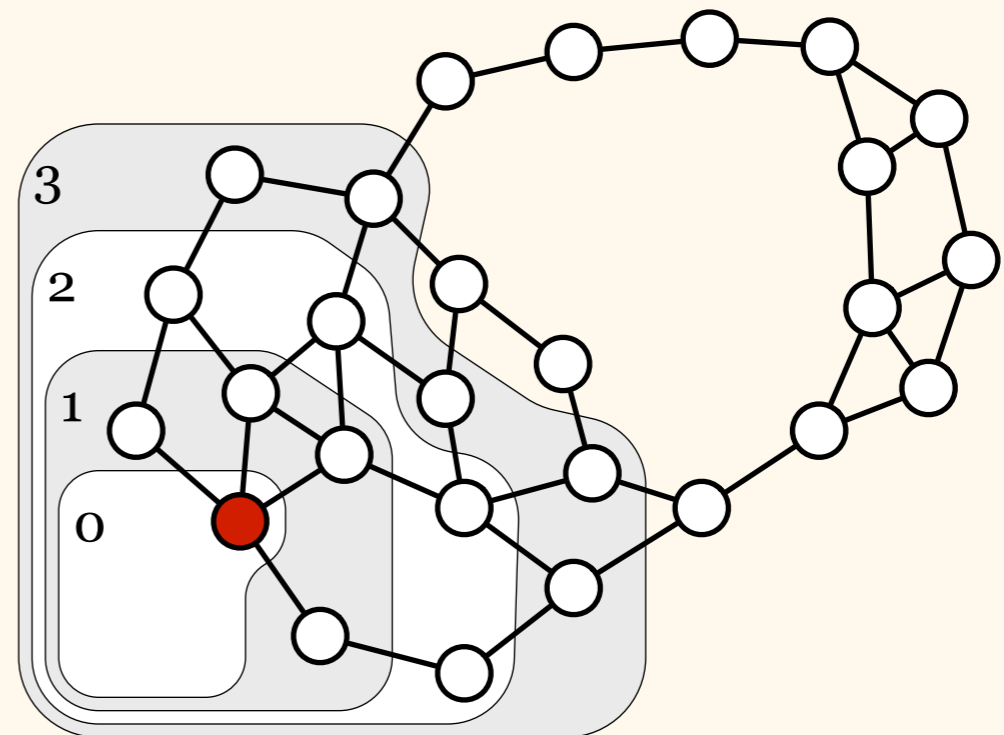
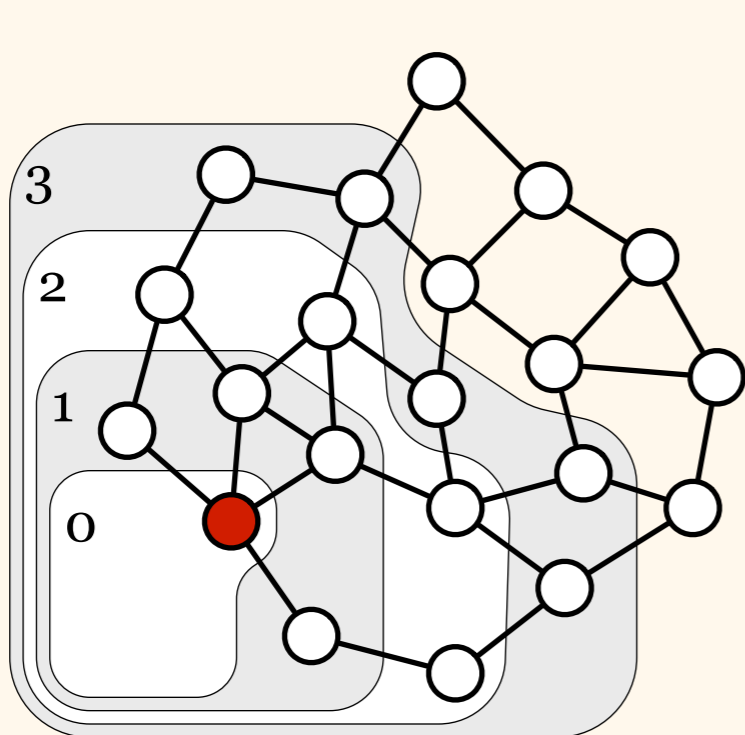


Local Neighbourhoods

- Local neighbourhoods of nodes u and v “look identical” up to distance r
 - isomorphism between radius- r neighbourhood of u and radius- r neighbourhood of v
 - preserves *inputs, degrees, connections, and port numbers*

Local Neighbourhoods

- Local neighbourhoods of nodes u and v “look identical” up to distance r

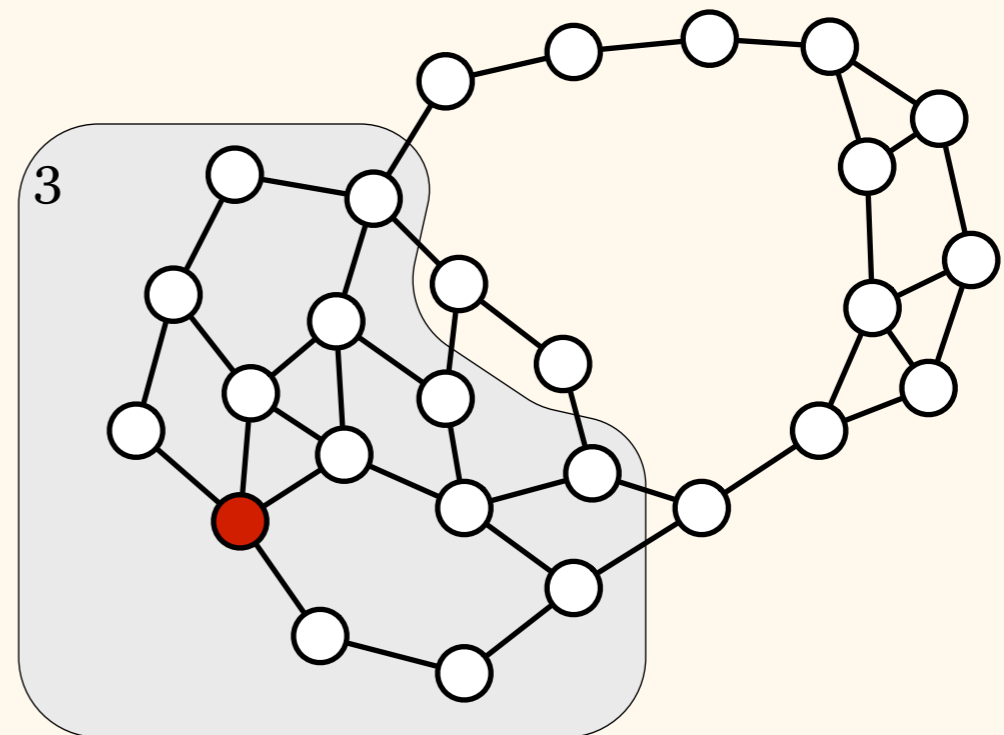
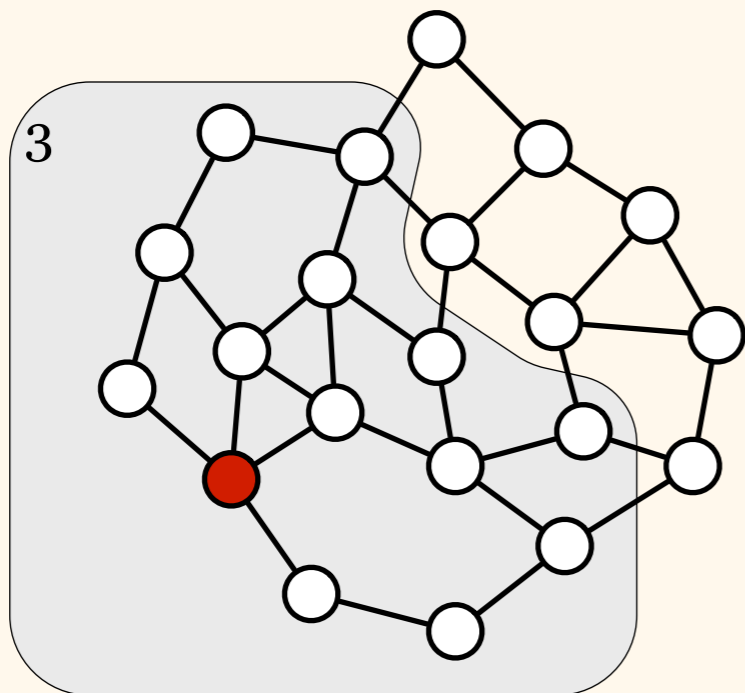


Local Neighbourhoods

- Local neighbourhoods of nodes u and v “look identical” up to distance r
- **Theorem:** In any algorithm, up to time r , the local states of u and v are identical
- *Informal proof:* time \approx distance
- *Formal proof:* by induction on time

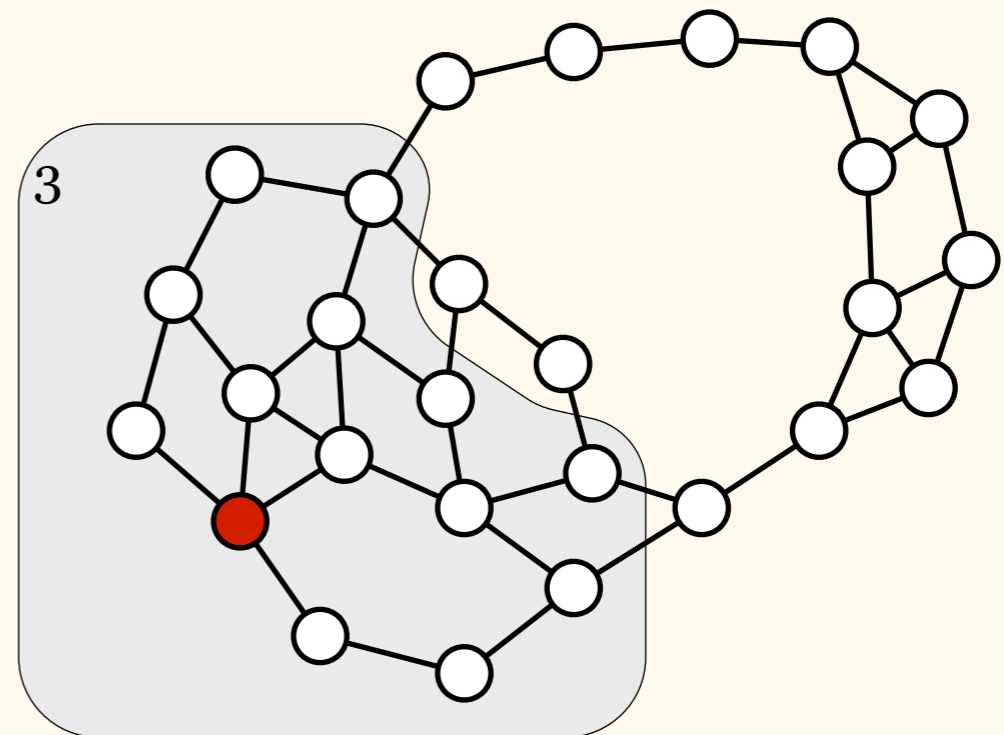
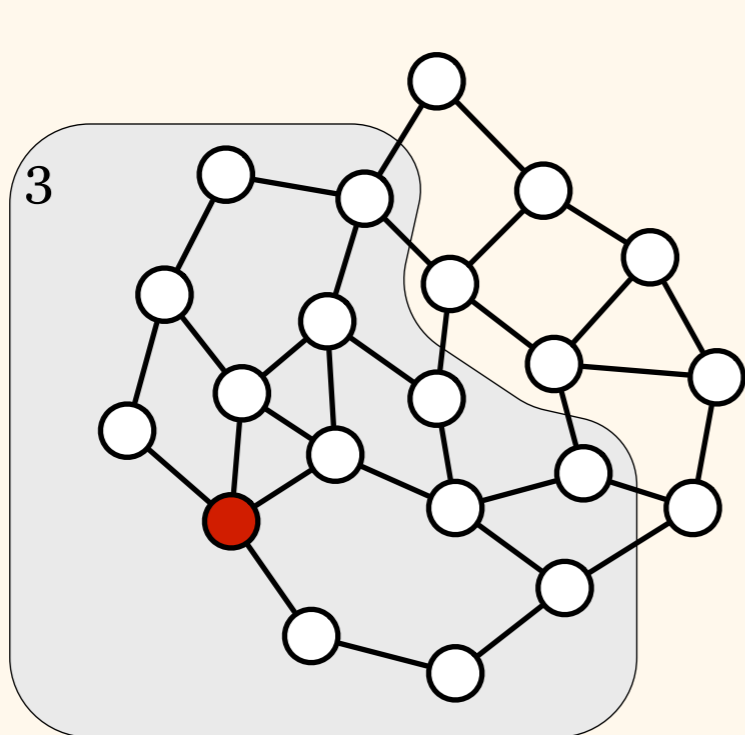
Local Neighbourhoods

- Time 0: identical *local states* in radius- r neighbourhoods



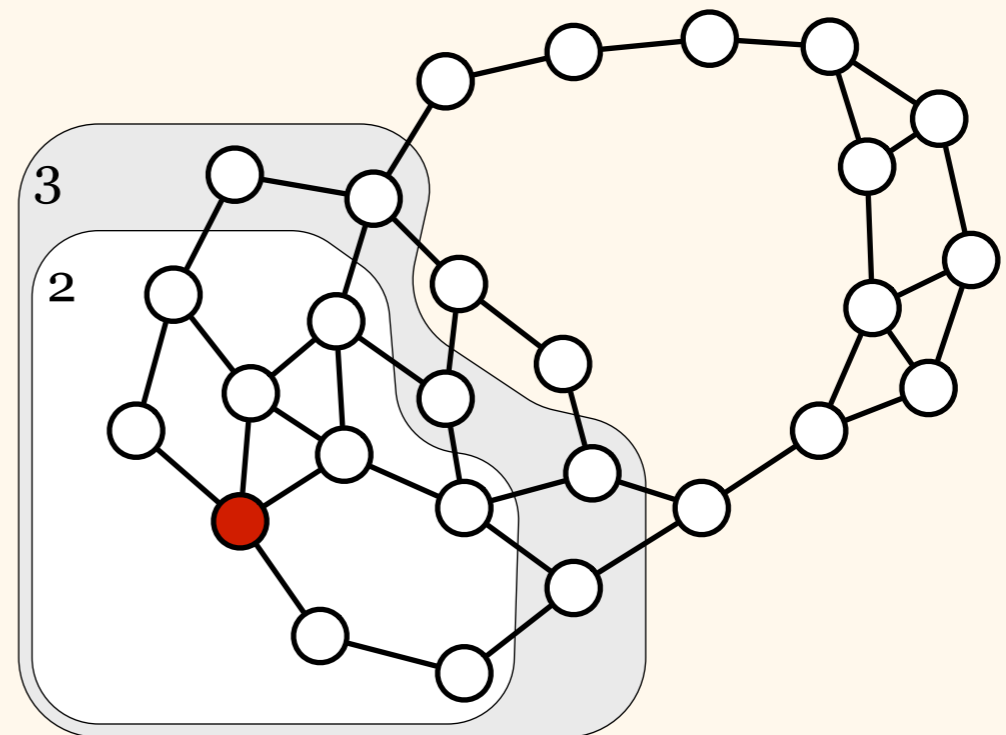
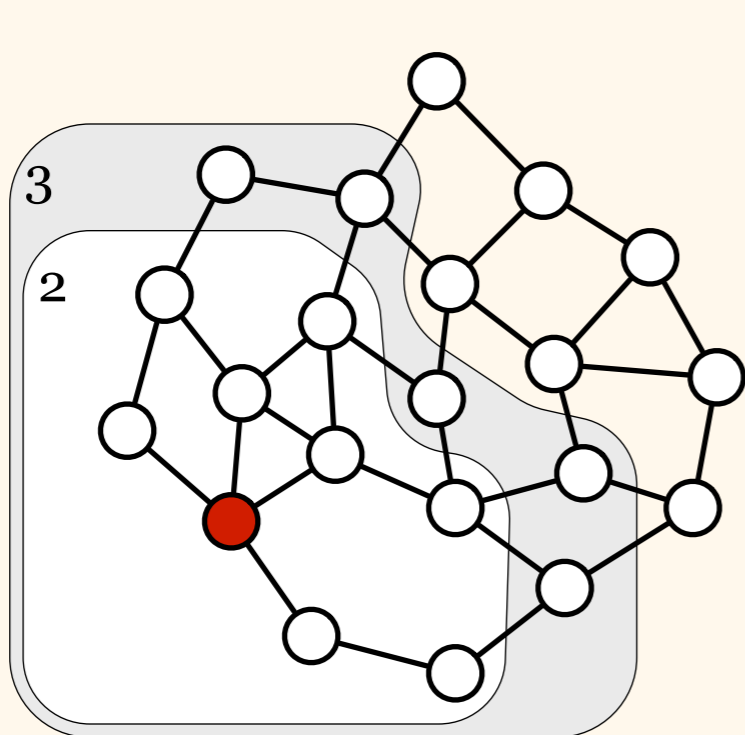
Local Neighbourhoods

- Time 1: identical *outgoing messages* in radius- r neighbourhoods



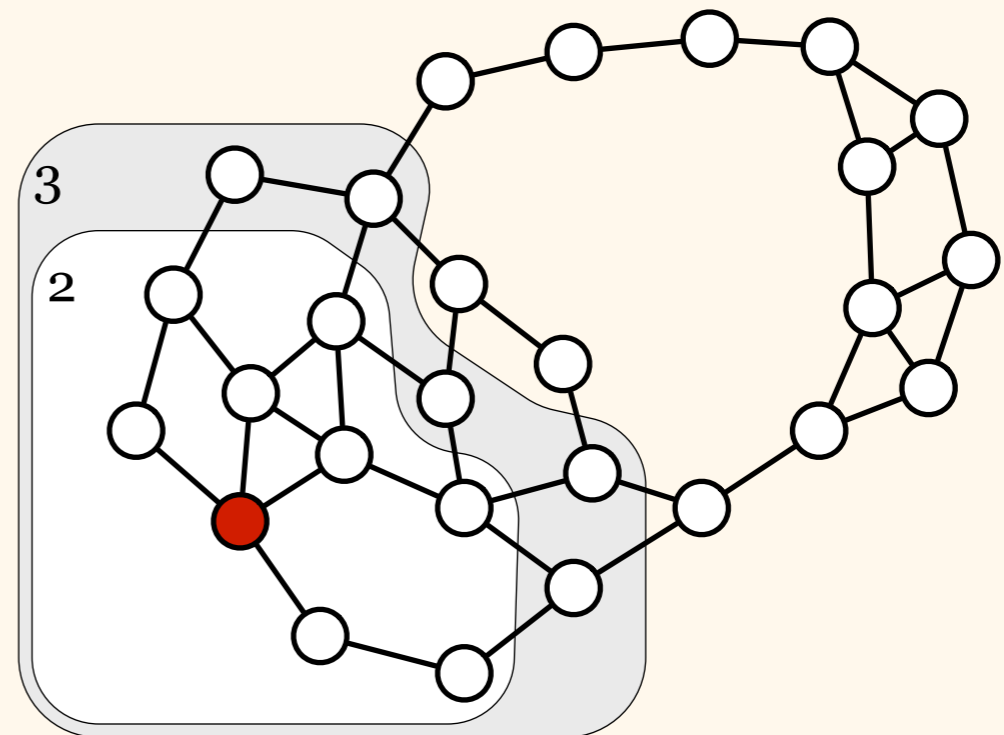
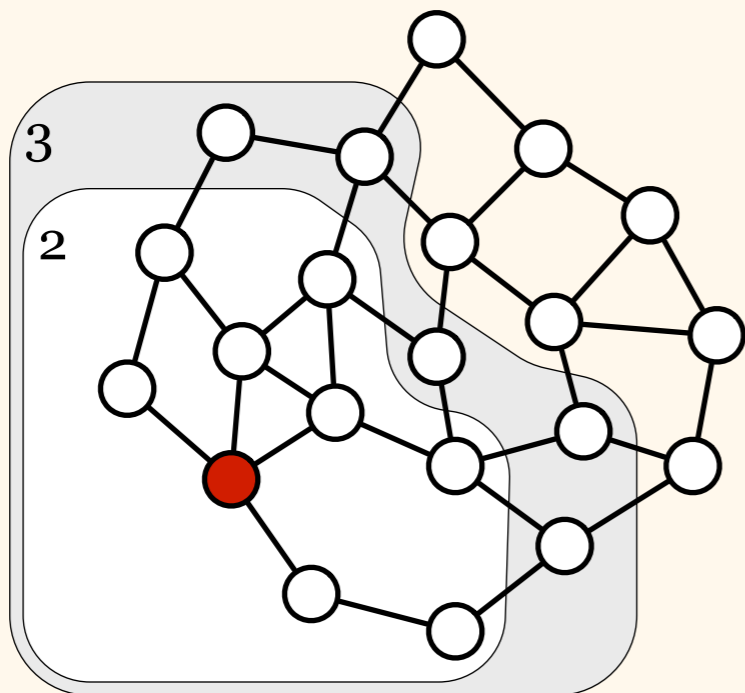
Local Neighbourhoods

- Time 1: identical *incoming messages* in radius- $(r-1)$ neighbourhoods



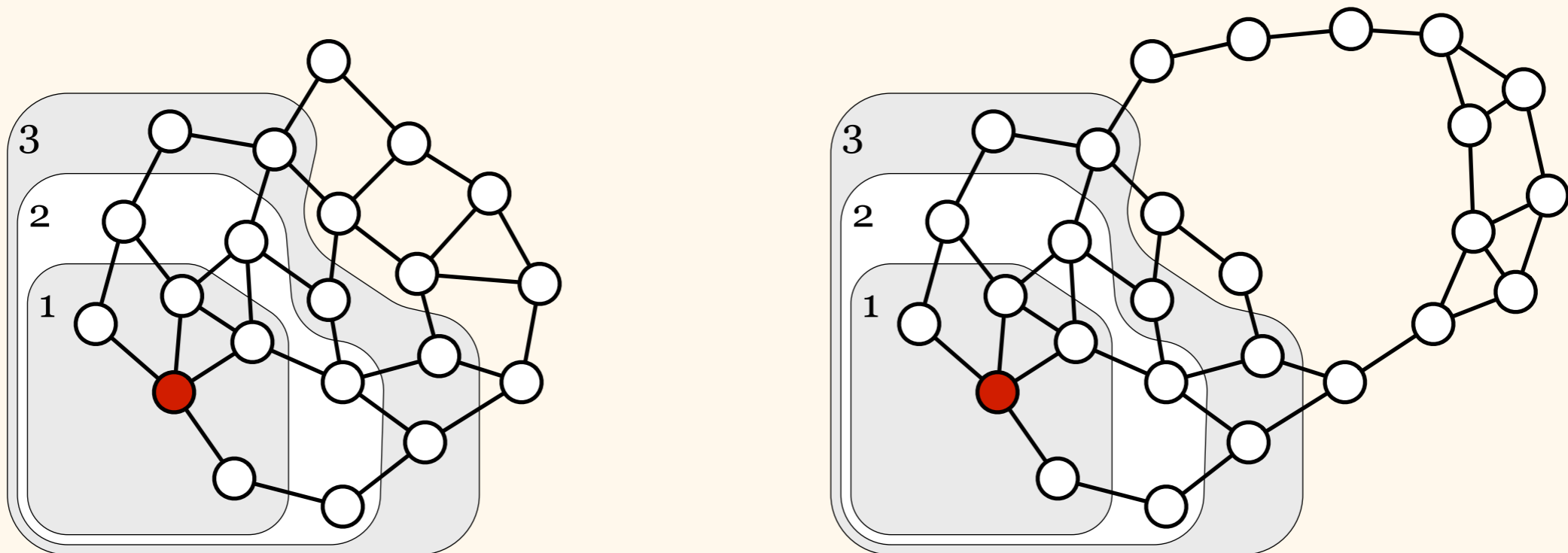
Local Neighbourhoods

- Time 1: identical *local states* in radius- $(r-1)$ neighbourhoods



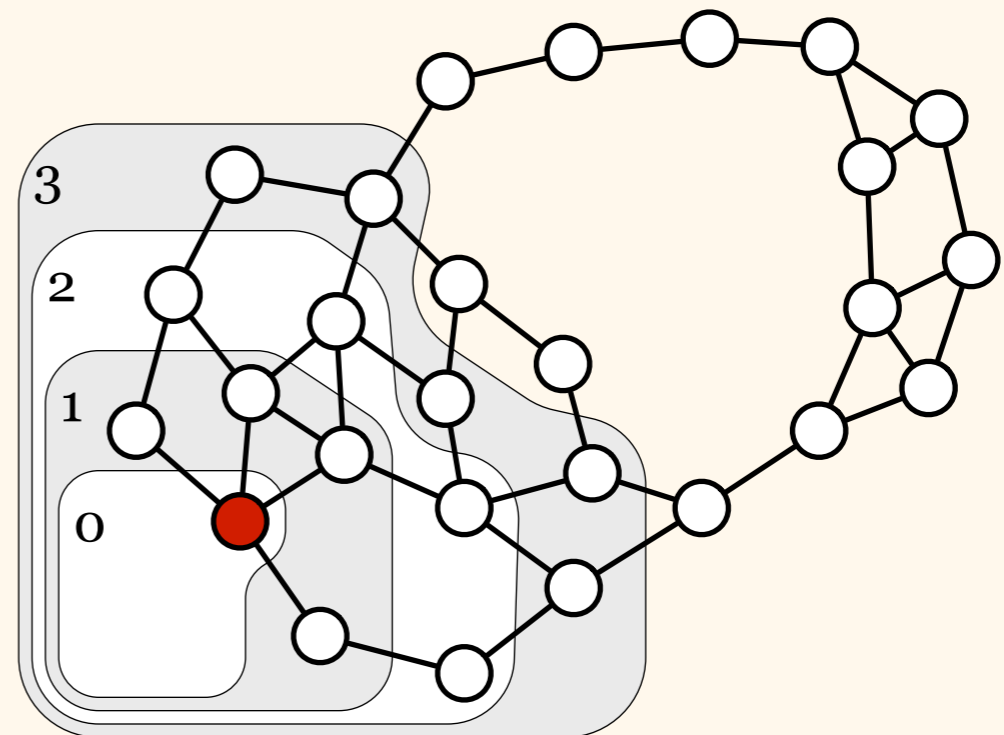
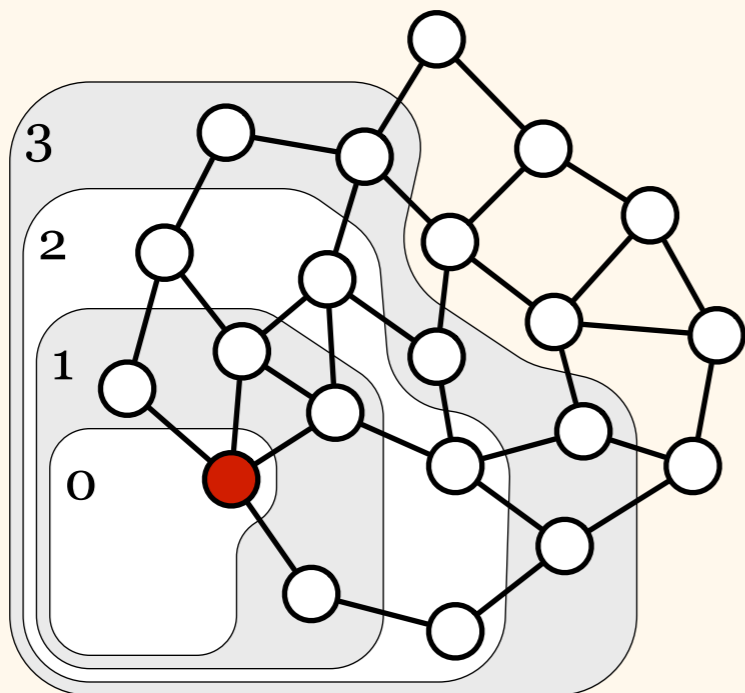
Local Neighbourhoods

- Time t : identical *local states* in radius- $(r-t)$ neighbourhoods



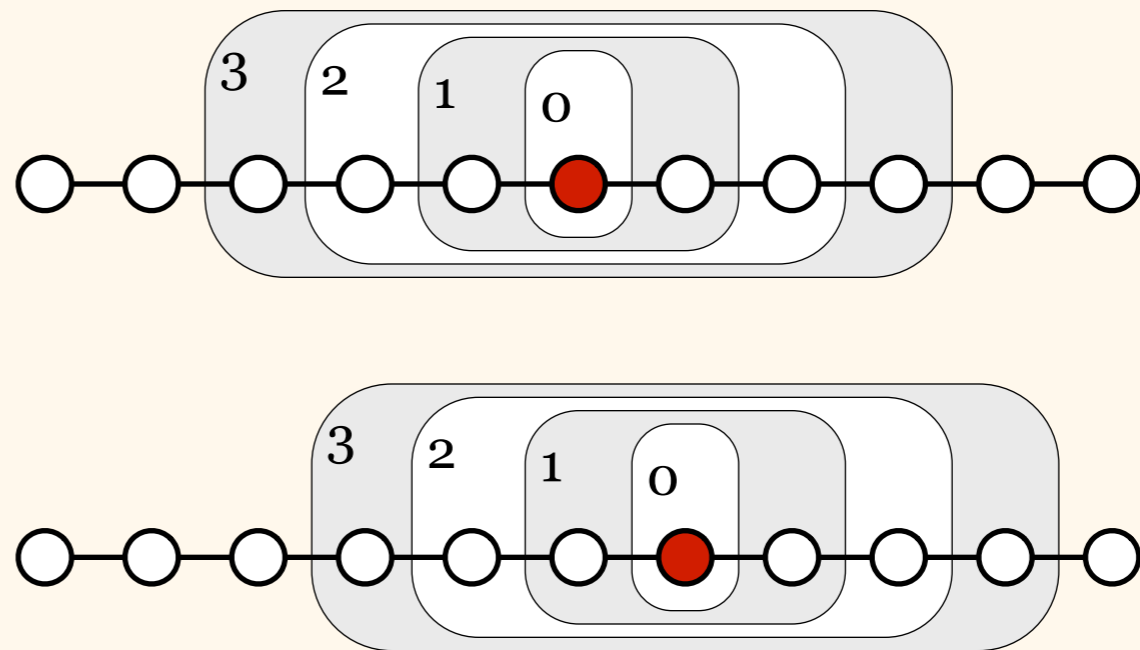
Local Neighbourhoods

- Time r : identical *local states* in radius-0 neighbourhoods



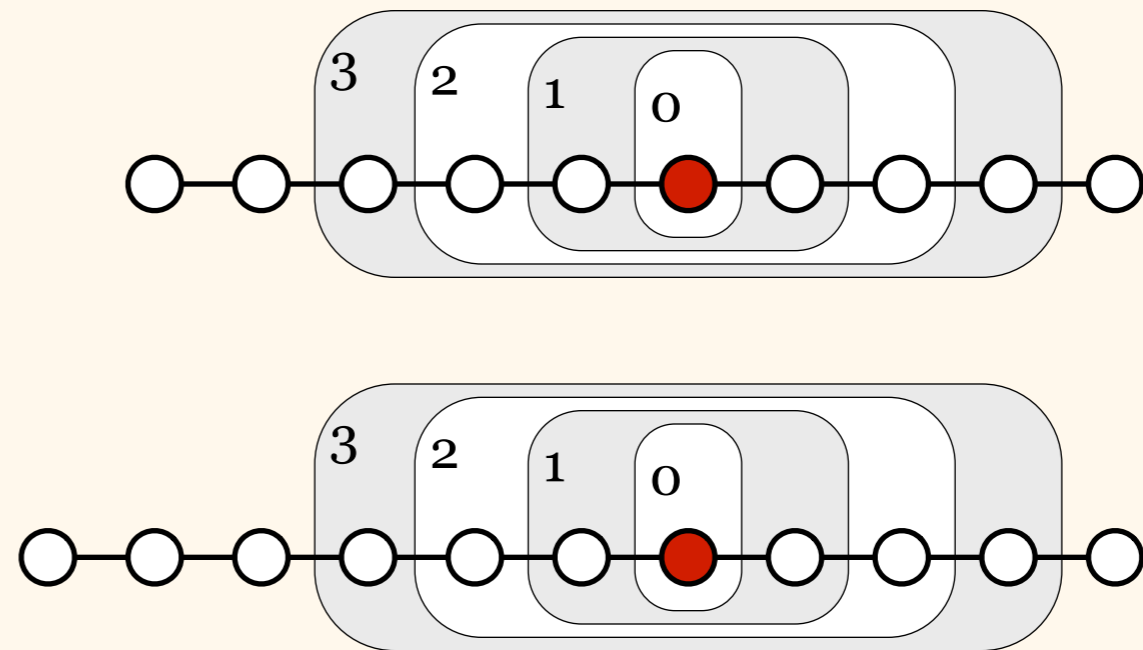
Local Neighbourhoods

- Application: finding midpoint of a path requires $\Omega(n)$ rounds



Local Neighbourhoods

- Application: counting the number of nodes requires $\Omega(n)$ rounds



Proof Techniques

- Covering maps
 - problems that cannot be solved at all
- Isomorphic local neighbourhoods
 - problems that cannot be solved quickly
- Plenty of exercises...