

- **Weeks 1–2: informal introduction**

- network = path



- **Week 3: graph theory**

- **Weeks 4–7: models of computing**

- what can be computed (efficiently)?

- **Weeks 8–11: lower bounds**

- what cannot be computed (efficiently)?

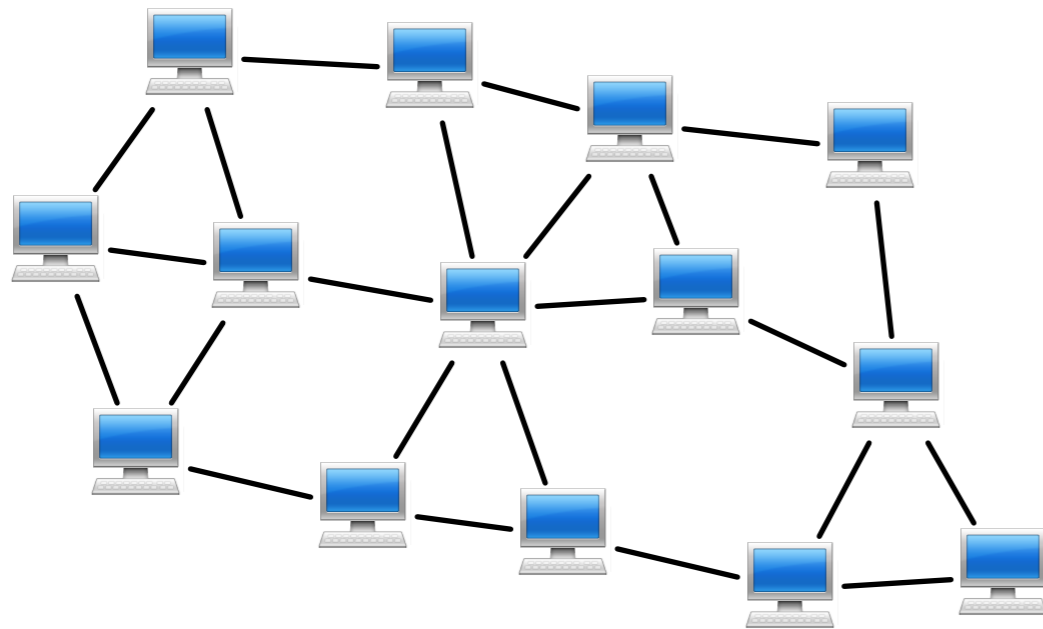
- **Week 12: recap**

Week 12

– Conclusions

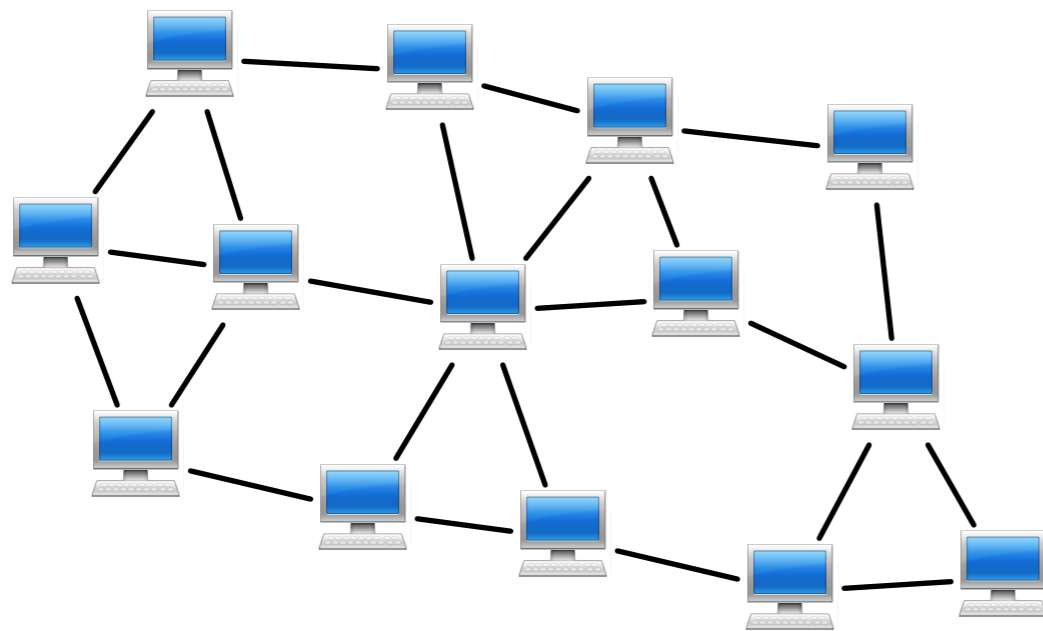
Recap: **Distributed algorithms**

Algorithms for computer networks



Recap: **Distributed algorithms**

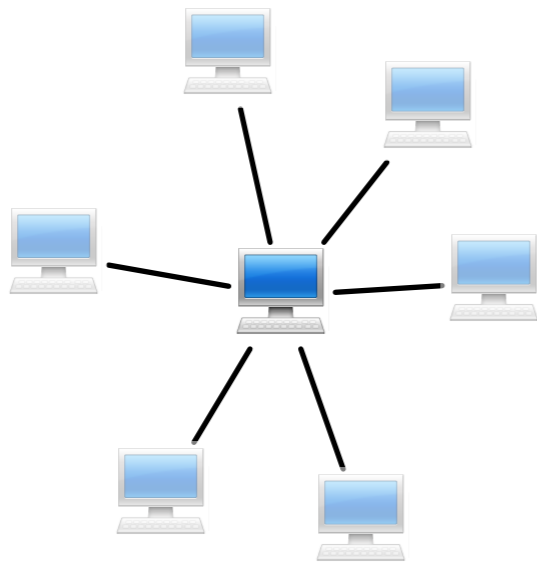
**Identical computers in an unknown network,
all running the same algorithm**



Recap:

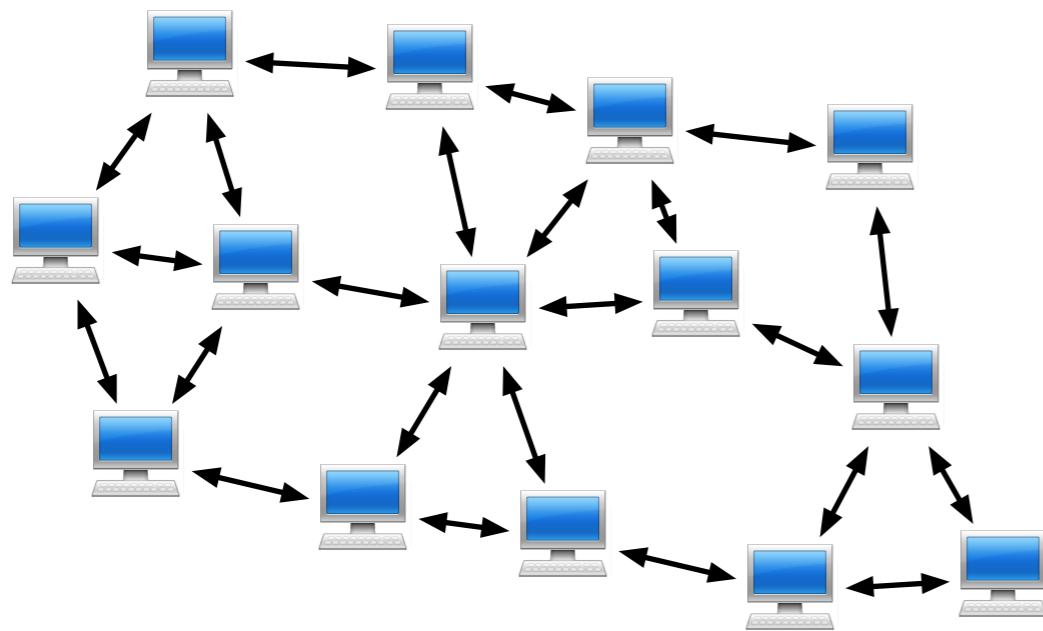
Distributed algorithms

Initially each computer only aware of its immediate neighbourhood



Recap: **Distributed algorithms**

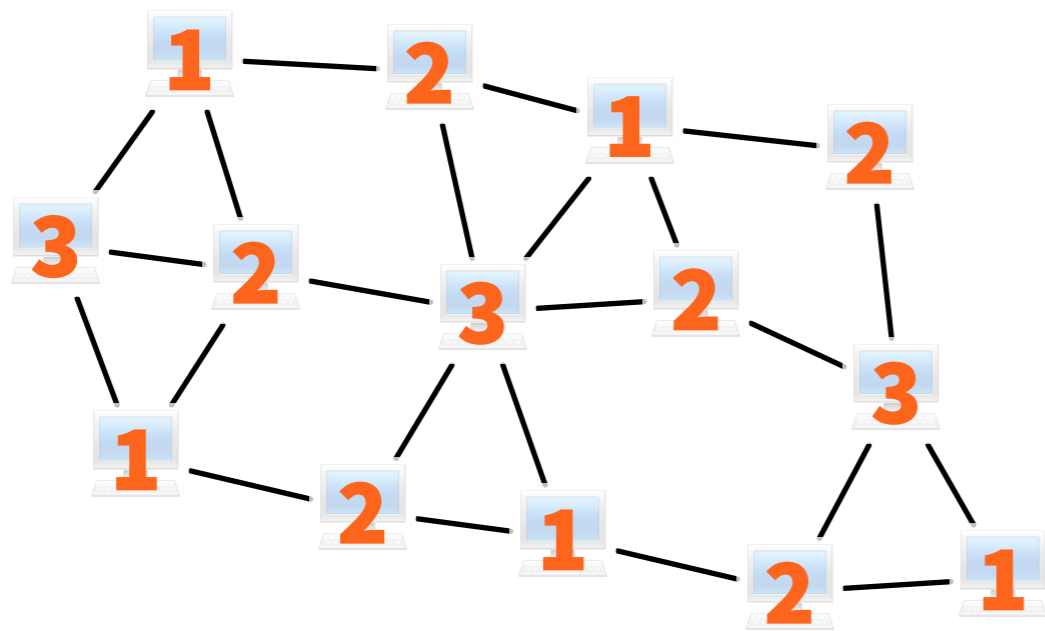
**Nodes can exchange messages
with their neighbours to learn more...**



Recap:

Distributed algorithms

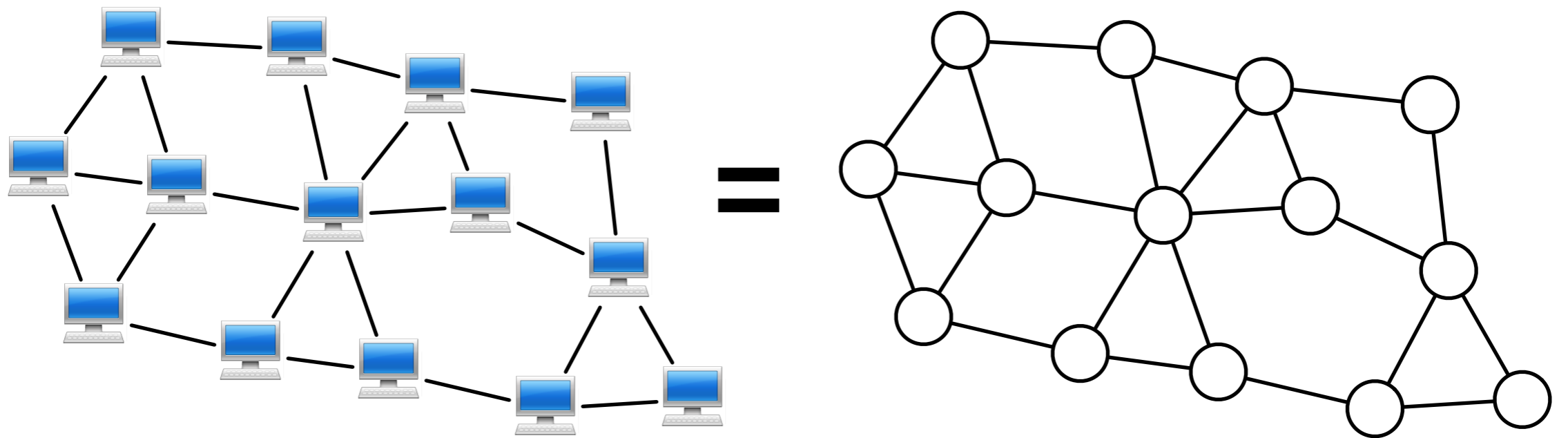
Finally, each computer has to stop and produce its own **local output**



Recap:

Distributed algorithms

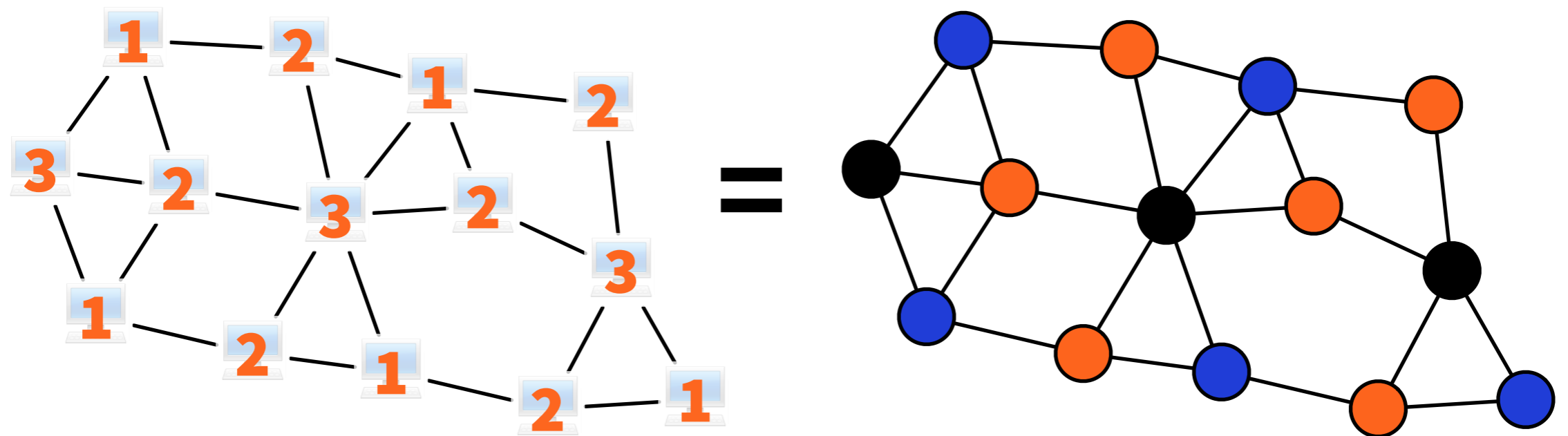
Focus on graph problems:
network topology = input graph



Recap: **Distributed algorithms**

Focus on graph problems:

local outputs = solution (here: graph colouring)



Recap:

Distributed algorithms

Typical research question:

“How fast can we solve graph problem X?”

Time = number of communication rounds

What have we learned?

- Dealing with *unknown systems*
- Dealing with *partial information*
- Dealing with *parallelism*
- **Applications beyond distributed computing:**
fault tolerance, online, streaming, multicore...

Learning objectives

- **Models**
- **Algorithms**
- **Lower bounds**
- **Graph theory**

Objective 1:

Models of computing

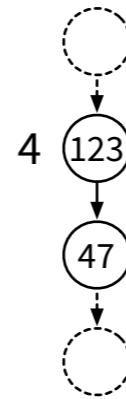
- **Precisely** what is a “*distributed algorithm*”
- **In each of these models:**
 - PN, LOCAL, CONGEST
 - deterministic, randomised

Objective 2: **Algorithms**

- **Colouring paths:** LOCAL, $O(\log^* n)$
- **Colouring graphs:** LOCAL, $O(\log n)$ w.h.p.
- **Gather everything:** LOCAL, $O(\text{diam}(G))$
- **Bipartite maximal matching:** PN, $O(\Delta)$
- **All-pairs shortest paths:** CONGEST, $O(n)$

Algorithm P3CBit: Fast colour reduction

$c_0 = 123 = 01111011_2$ (my colour)
 $c_1 = 47 = 00101111_2$ (successor's colour)
 $i = 2$ (bits numbered 0, 1, 2, ... from right)
 $b = 0$ (in my colour bit number i was 0)
 $c = 2 \cdot 2 + 0 = 4$ (my new colour)



$k = 8$, reducing from $2^8 = 256$ to $2 \cdot 8 = 16$ colours

Algorithm P3CBit.

Fas

$c_0 = 12$

$c_1 = 4$

$i = 2$ (

$b = 0$ (

$c = 2 \cdot 2$

k

Algorithm idea 3

- Colour palette: $\{1, 2, \dots, \Delta + 1\}$
- Active with probability $1/2$
- If *active*, pick a random *free* colour
 - not used by any neighbour that has stopped
- Try again if conflicts...

Algorithm P3CBit.

Fast

$c_0 = 12$

$c_1 = 4$

$i = 2$

$b = 0$

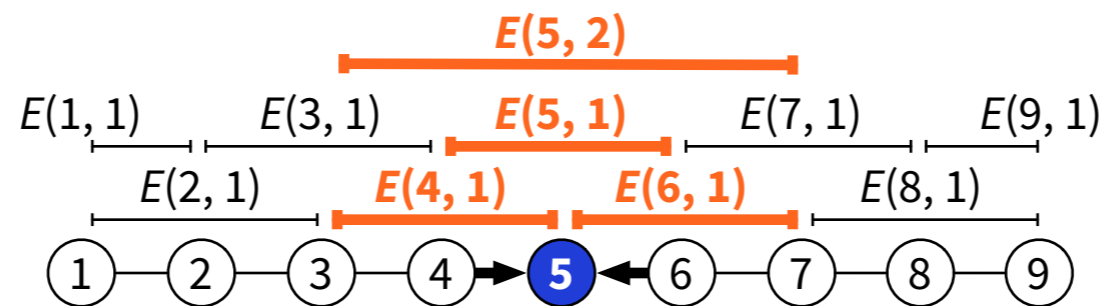
$c = 2 \cdot 2$

Algo

- Color
- Activ
- If act
- no
- Try a

Gathering everything

- Given $E(v, r)$, we can learn $E(v, r + 1)$ in 1 round
 - send $E(v, r)$ to all neighbours, take union



Algorithm P3CBit.

Fast

$c_0 = 12$

$c_1 = 4$

$i = 2$

$b = 0$

$c = 2 \cdot 2$

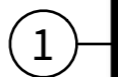
Algo

- Color
- Active
- If *act*
- Try a

Gat

- Given
- se

$E(1, 1)$



Algorithm BMM: Maximal matching

- **Blue nodes** send proposals to their orange neighbours one by one
 - using port numbers
- **Orange nodes** accept the first proposal that they get
 - using port numbers to break ties

Algorithm P3CBit.

Fast

$c_0 = 12$

$c_1 = 4$

$i = 2$

$b = 0$

$c = 2 \cdot 2$

Algo

- Color
- Active
- If *act*
- Try a

Gate

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- se

$E(1, 1)$

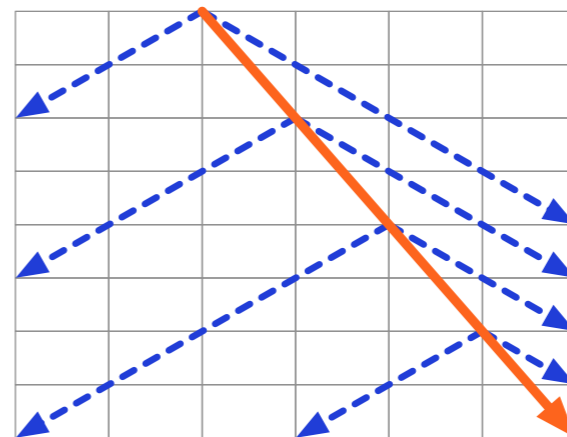
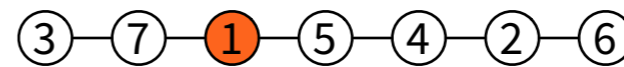


Algorithm BMM.

Max

- Blue
- orange
- US
- Oran
- the f
- US

Algorithm APSP



wave

token

Objective 2: **Algorithms**

- **Reductions!**
- **Graph colouring is a very useful subroutine**

Objective 3:

Lower bounds

- **Covering maps:**
what cannot be solved at all in PN model
- **Local neighbourhoods:**
what cannot be solved fast in any model
- **Ramsey's theorem:**
what cannot be solved in $O(1)$ time

Objective 4:

Graph theory

- **Basic definitions**
- **Connections between graph problems**
 - e.g. maximal matching \rightarrow small vertex covers
- **Ramsey's theorem**
 - at least for $c = 2, k = 2$

What else is studied in distributed computing?

- Fault-tolerance
- Asynchrony
- Shared memory
- Physical models
- Robot navigation
- Nondeterminism
- Complexity measures
- High-performance computing
- Practical aspects of networking ...

What next?

- **ICS-E4020 Programming Parallel Computers**
 - 5th period, 5 credits, intensive course
 - programming modern parallel computers: multicore, GPU, memory hierarchies ...
 - hands-on programming exercises
 - main goal: make it as fast as you can!

What next?

- **Just ask if you want to do more!**
 - master's thesis topics?
 - summer internships?
 - doctoral studies?

Practicalities

- **2nd mid-term exam: 10 December**
 - remember to register on time!
- **Course feedback: deadline 17 December**
 - 1 extra point in grading

What to expect in the exam?

- **See the learning objectives!**
- **Do not think that you can safely forget what we learned during the 1st period!**
- **Expect both algorithm design and lower bound proofs**

Examples of old exam problems

- Prove: no deterministic PN-algorithm that finds a *minimum vertex cover* in cycle graphs, given a *minimal vertex cover*

Examples of old exam problems

- **Prove: no deterministic PN-algorithm that finds a 6-colouring in cycle graphs given a maximal independent set**

Examples of old exam problems

- **Counting problem:** all nodes output $|V|$
- **Prove: no deterministic PN-algorithm for cycle graphs**
- **Prove: no $o(n)$ -time deterministic LOCAL-algorithm for cycle graphs**

Examples of old exam problems

- **Prove: no deterministic PN-algorithm for maximal matching in arbitrary graphs**

Examples of old exam problems

- Prove: no deterministic $o(n)$ -time PN-algorithm for **weak 2-colouring** in paths of length ≥ 3

Examples of old exam problems

- Give an **elementary proof** that any graph with 6 nodes contains a clique with 3 nodes or an independent set with 3 nodes

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- **Week 12: recap**