- Weeks 1–2: informal introduction
 - network = path



- Week 3: graph theory
- Weeks 4–7: models of computing
 - what can be computed (efficiently)?
- Weeks 8–11: lower bounds
 - what cannot be computed (efficiently)?
- Week 12: recap

Mid-term exams

- Mid-term exams:
 - Thursday, 23 October 2014, 9:00am
 - Thursday, 11 December 2014, 9:00am
- Register on time (one week before) in Oodi

Mid-term exams

- Topics:
 - 1st exam: Chapters 1–6
 - 2nd exam: Chapters 1–12
- See course web page for learning objectives!

Week 6

 CONGEST model: bandwidth limitations

- LOCAL model: arbitrarily large messages
- CONGEST model: O(log n)-bit messages

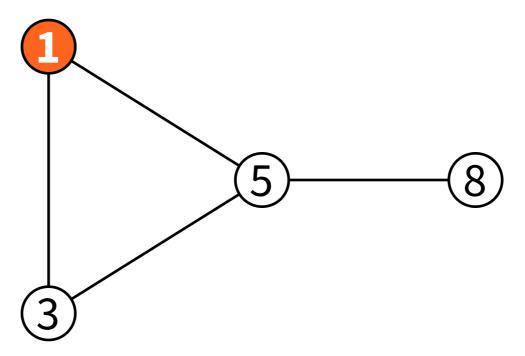
- Any of these can be encoded in O(log n)-bit messages:
 - node identifier
 - number of nodes
 - number of edges
 - distance between two nodes ...

- Many algorithms that we have seen only send small messages
 - can be used directly in the CONGEST model
- Exception: algorithm Gather
 - may need to send $O(n^2)$ -bit messages

- O(n) time trivial in the LOCAL model
 - brute force approach: Gather + solve locally
- O(n) time non-trivial in the CONGEST model
- Today: how to find all-pairs shortest paths in O(n) time

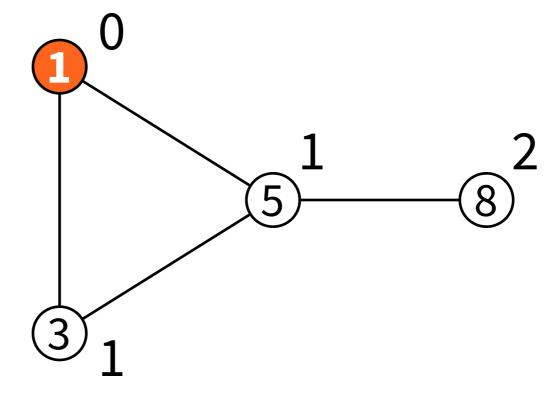
Single-source shortest paths

Input:



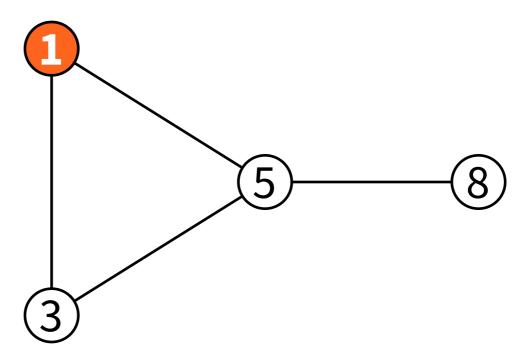
Single-source shortest paths

Output:



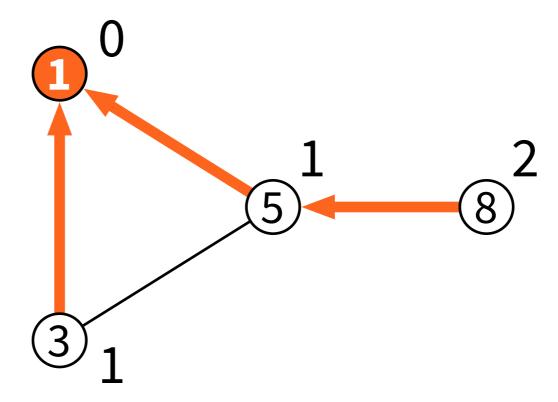
BFS tree

Input:



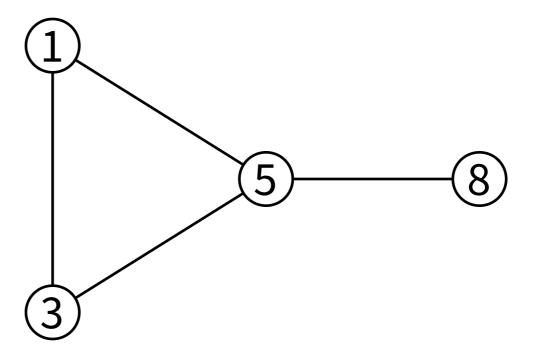
BFS tree

Output:



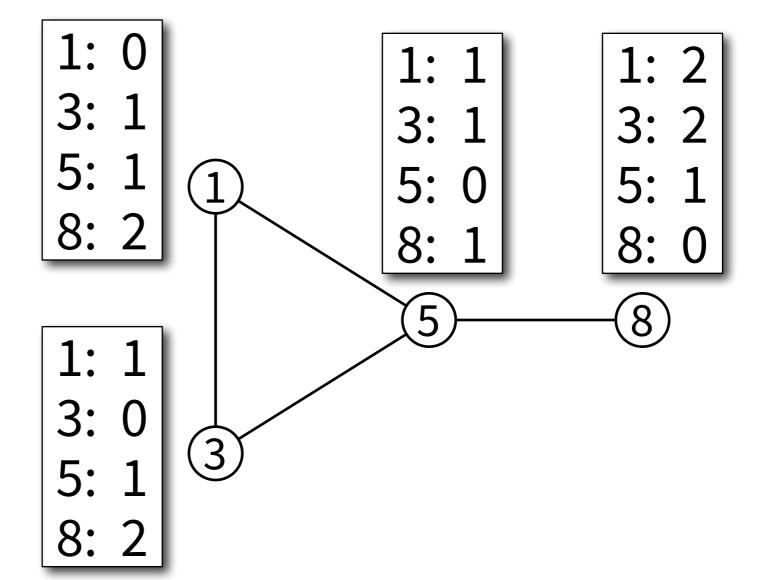
All-pairs shortest paths

Input:



All-pairs shortest paths

Output:

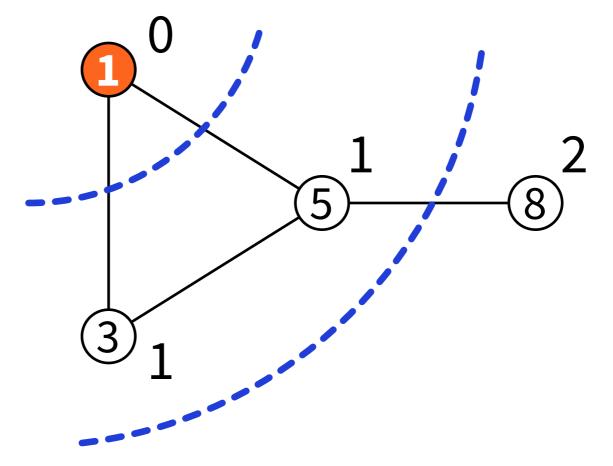


Algorithm Wave

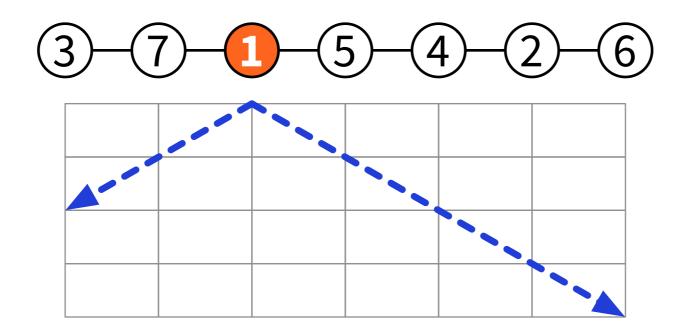
- Solves single-source shortest paths in time O(diam(G))
- Leader creates a 'wave', other nodes propagate it
- Wave first received in round t: distance to leader is t

Algorithm Wave

Output:



Algorithm Wave

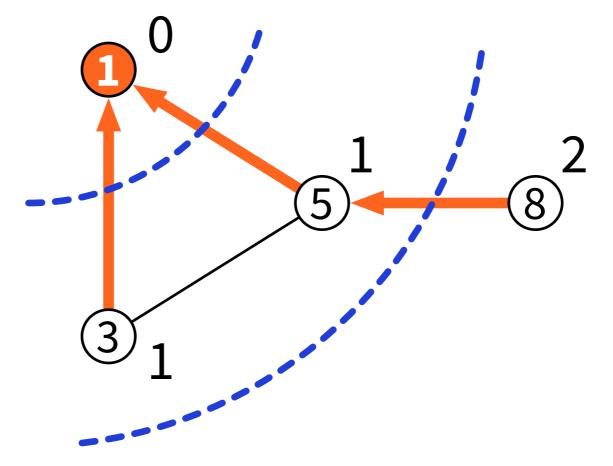


Algorithm BFS

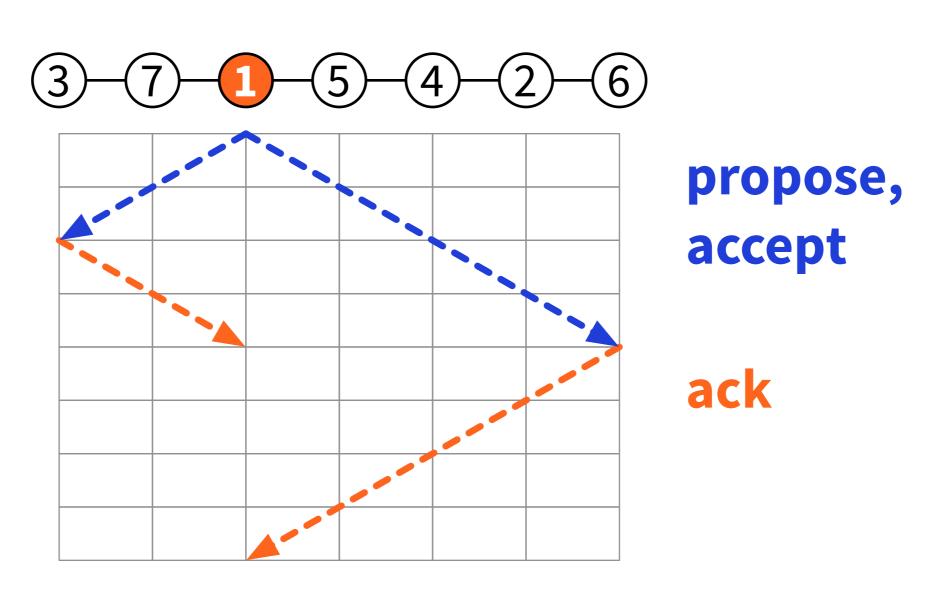
- Wave + handshakes
- Tree construction:
 - "proposal" + "accept"
 - everyone knows their parent & children
- Acknowledgements back from leaf nodes

Algorithm BFS

Output:



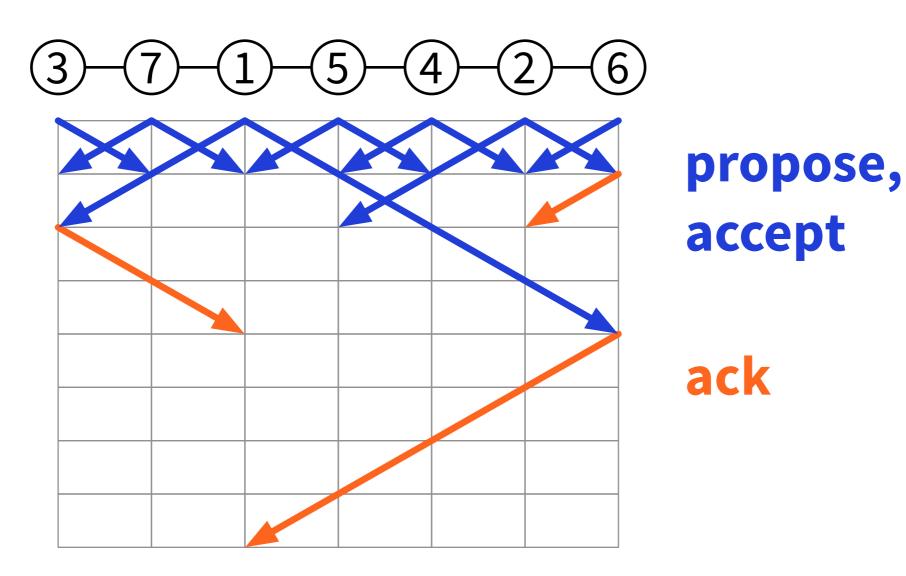
Algorithm BFS



Algorithm Leader

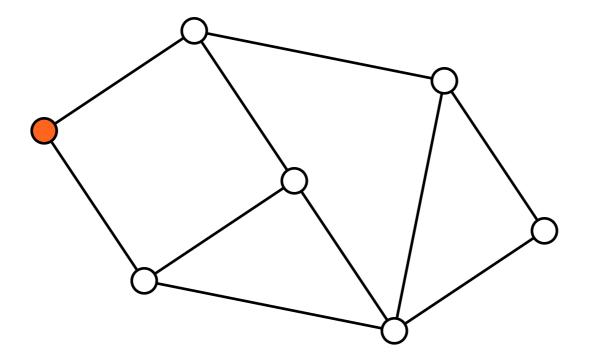
- Each node creates a separate BFS process
- When two BFS processes "collide", the one with the smaller root "wins"
 - each node only needs to send messages related to one BFS process
- One tree wins everyone else → leader

Algorithm Leader

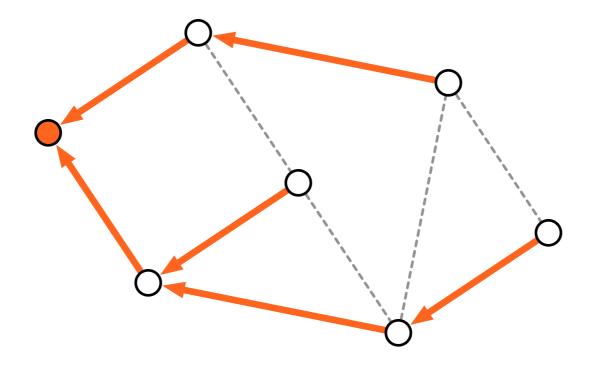


- Basic idea: run Wave from each node
- Challenge: congestion
 - all waves parallel → too many bits per edge
 - all waves sequentially → takes too long
- Solution: pipelining

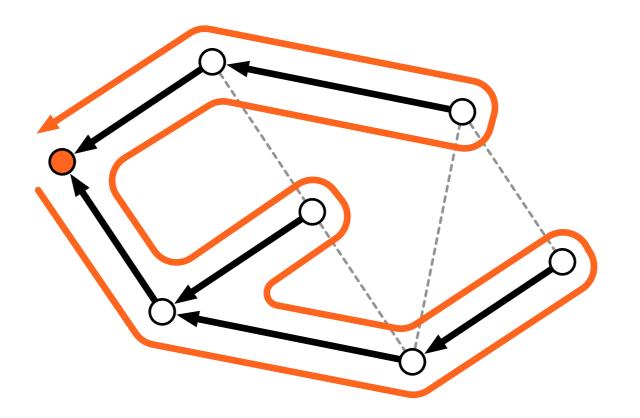
Elect leader



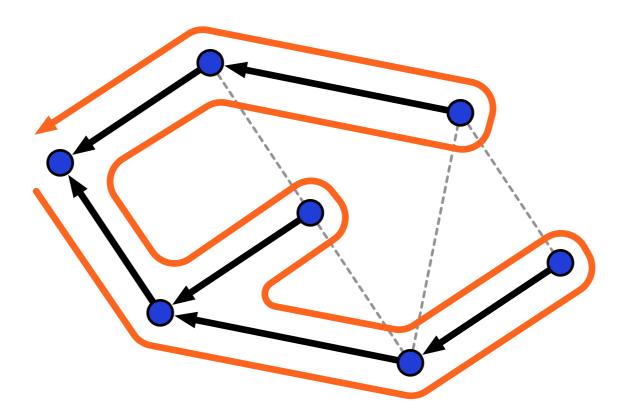
Elect leader, construct BFS tree

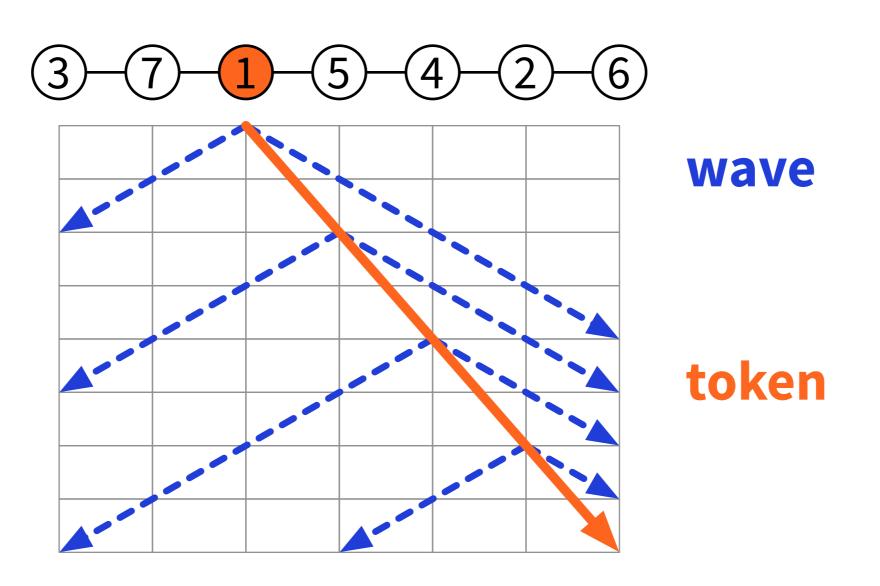


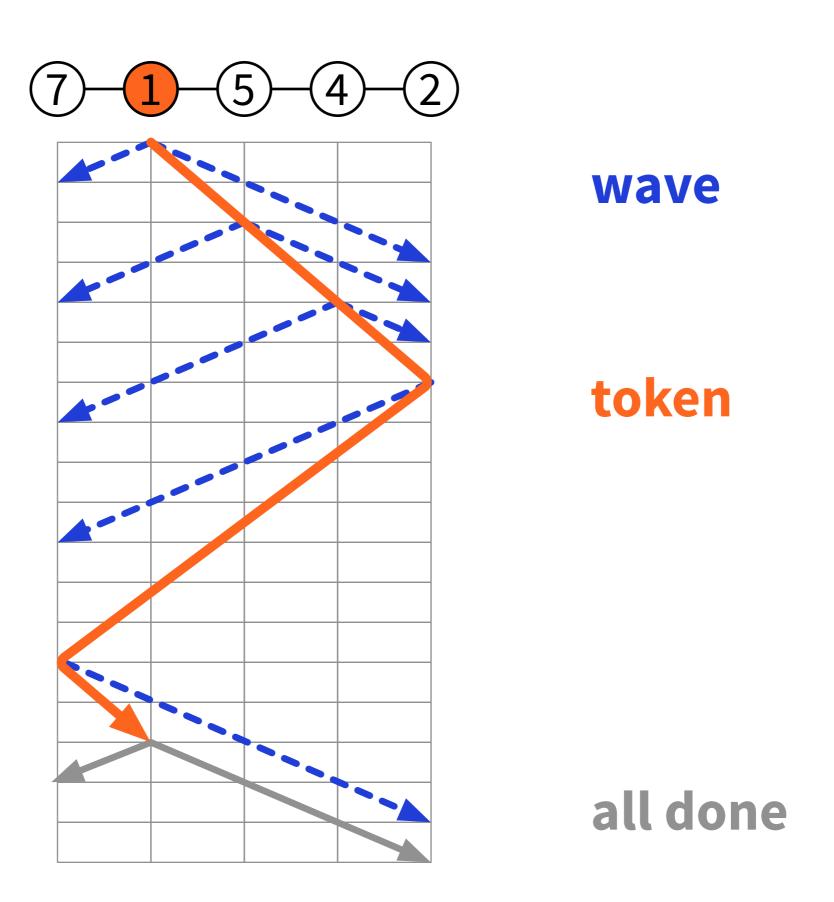
Move token along BFS tree slowly



Create wave every time we visit a new node







Algorithm animation:

http://users.ics.aalto.fi/suomela/apsp/

Pipelining

- n operations, each operation takes time t
- Parallel: t rounds, bad congestion
- Sequential: nt rounds, no congestion
- Pipelining: n + t rounds, no congestion

Summary

- LOCAL model: unlimited bandwidth
- CONGEST model: O(log n) bandwidth
- O(n) or O(diam(G)) time is no longer trivial
- Example: all-pairs shortest paths in time O(n), pipelining helps

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